



ORTHOGONAL COLOURINGS OF TENSOR GRAPHS

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ABSTRACT. Perfect k -orthogonal colourings of tensor product graphs are studied in this article. First, the problem of determining if a given graph has a perfect 2-orthogonal colouring is reformulated as a tensor subgraph problem. Then, it is shown that if two graphs have a perfect k -orthogonal colouring, then so does their tensor graph. This provides an upper bound on the k -orthogonal chromatic number for general tensor graphs. Lastly, two other conditions for a tensor graph to have a perfect k -orthogonal colouring are given.

1. INTRODUCTION

Two proper colourings of a graph are *orthogonal* if when two elements are coloured with the same colour in one of the colourings, then those elements receive distinct colours in the other colouring. Archdeacon, Dinitz, and Harary [2] originally studied this type of colouring, in the context of edge colourings. Then, Caro and Yuster [4] revisited this concept, this time in the context of vertex colouring. In this paper, the vertex variation is studied.

A k -orthogonal colouring of a graph G is a collection of k mutually orthogonal vertex colourings. For simplicity, a 2-orthogonal colouring is called an *orthogonal colouring*. The k -orthogonal chromatic number of a graph G , denoted by $O\chi_k(G)$, is the minimum number of colours required for a proper k -orthogonal colouring. Again for simplicity, the 2-orthogonal chromatic number is simply denoted by $O\chi(G)$ and simply called the *orthogonal chromatic number*.

For a graph G with n vertices, $O\chi_k(G) \geq \lceil \sqrt[n]{n} \rceil$. Otherwise, there are fewer colour k -tuples than there are vertices. If G has n^2 vertices and $O\chi_k(G) = \lceil \sqrt[n]{n^2} \rceil = n$, then G is said to have a *perfect k -orthogonal colouring*. A perfect 2-orthogonal colouring is simply called a *perfect orthogonal colouring*. Perfect orthogonal colourings are of particular importance because they have applications to independent coverings [10] and scoring games [1].

Received by the editors January 14, 2022, and in revised form May 22, 2024.

2000 *Mathematics Subject Classification*. 05C15.

Key words and phrases. Graph Colouring, Graph Product, Graph Homomorphism.

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Due to these applications, research is focused on determining which graphs have perfect k -orthogonal colourings. For instance, Caro and Yuster [4] constructed graphs having perfect k -orthogonal colourings by using orthogonal Latin squares. Ballif [3] studied upper bounds on sets of orthogonal colourings. Whereas Janssen and the author [7] studied perfect k -orthogonal colourings of circulant graphs.

In this paper, the tensor graph product is used to construct graphs having perfect k -orthogonal colourings. The *tensor product* of two graphs G and H , denoted by $G \times H$, has vertex set $V(G) \times V(H)$, and two vertices (u_1, v_1) and (u_2, v_2) in $G \times H$ are adjacent if and only if $u_1 u_2 \in E(G)$ and $v_1 v_2 \in E(H)$. If G is a graph that is created by the tensor product of two graphs, then G is called a *tensor graph*.

For graph products, there can be a clear relationship between colourings of the factors and colourings of the product graph. For example, it is well-known that the chromatic number of a Cartesian graph (see Section 3 for a formal definition) is the minimum of the chromatic numbers of the factors. For tensor graphs, it was conjectured by Hedetniemi [5] that the same result would be true. However, Hedetniemi's conjecture was recently disproved by Shitov [11].

In this paper, it is shown that for a graph G , $O\chi(G) \leq n$ if and only if G is a subgraph of the tensor graph $K_n \times K_n$. Then, it is shown that if two graphs have a perfect k -orthogonal colouring, then so does their tensor graph. For $k = 2$, only one factor is required to have a perfect orthogonal colouring. Similarly, it is shown that this condition can be relaxed for perfect k -orthogonal colourings. These results were first obtained in the author's doctoral dissertation [9].

In the language of graph homomorphisms, a t -colouring of G is a homomorphism $f : G \rightarrow K_t$. A standard fact of homomorphisms is that there are homomorphisms ϕ and ψ such that $\phi : G \rightarrow X$ and $\psi : G \rightarrow Y$ if and only if $\phi \times \psi : G \rightarrow X \times Y$. In the context of this paper, we show in Theorem 2.1 that G admits an orthogonal t -colouring if and only if G admits an injective homomorphism $G \rightarrow K_t \times K_t$. We recommend the reader see [6] for more information on homomorphisms.

2. PERFECT ORTHOGONAL COLOURINGS

In this section, perfect orthogonal colourings of tensor graphs are studied. To start, it is shown that a graph G has $O\chi(G) \leq n$ if and only if is a subgraph of $K_n \times K_n$, which in this paper, is denoted by $G \subseteq K_n \times K_n$. Therefore, a graph G with m^2 vertices has a perfect orthogonal colouring if and only if it is a subgraph of $K_m \times K_m$.

Theorem 2.1. *For a graph G , $O\chi(G) \leq n$ if and only if $G \subseteq K_n \times K_n$.*

Proof. For $1 \leq i, j \leq n$, let (i, j) denote the vertices of the graph $K_n \times K_n$. First, suppose that $G \subseteq K_n \times K_n$. It will be shown that $K_n \times K_n$ has a perfect orthogonal colouring. If this is the case, then the orthogonal colouring of

$K_n \times K_n$ restricted to G is an orthogonal colouring of G using n colours, giving $O\chi(G) \leq n$.

Assign the vertex (i, j) the colour i in the first colouring and the colour j in the second colouring. For example, this orthogonal colouring is applied to $K_3 \times K_3$ in Figure 1. Displayed next to each vertex are the colours assigned in the first and second colouring.

Note that this assignment of colours has no orthogonal conflicts. It remains to check that there are no colour conflicts. Now, by the definition of the tensor product, for $1 \leq i_1, i_2, j_1, j_2 \leq n$, two vertices (i_1, j_1) and (i_2, j_2) in $K_n \times K_n$ are adjacent if and only if $i_1 \neq i_2$ and $j_1 \neq j_2$. Therefore, there are also no colour conflicts.

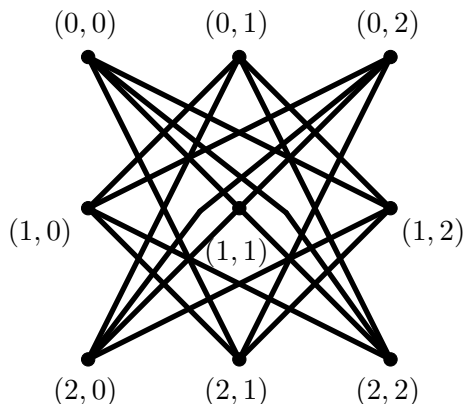


FIGURE 1. Orthogonal colouring of $K_3 \times K_3$

Now, suppose that $O\chi(G) \leq n$ and that (g_1, g_2) is an orthogonal colouring of G using the colours $\{1, 2, \dots, n\}$. To show that $G \subseteq K_n \times K_n$, an injective map that preserves edges is required. Let $F : G \rightarrow K_n \times K_n$ by $F(v) = (g_1(v), g_2(v))$. It is now shown that F is injective and preserves edges.

Since (g_1, g_2) is an orthogonal colouring of G , each colour pair is only assigned once. Thus, F is injective. Now, if $v_1 v_2 \in E(G)$, then $g_1(v_1) \neq g_1(v_2)$ and $g_2(v_1) \neq g_2(v_2)$ because g_1 and g_2 are proper. Therefore, we have that $(g_1(v_1), g_2(v_1))(g_1(v_2), g_2(v_2)) \in E(K_n \times K_n)$ by the definition of the edges in $K_n \times K_n$. Thus, F preserves edges. Since F is injective and preserves edges, $G \subseteq K_n \times K_n$. $\square$

Theorem 2.1 gives a way to reformulate the problem of determining if a graph has a perfect orthogonal colouring. This will be used later with the following theorem to obtain an upper bound on the orthogonal chromatic number of general tensor graphs. The following theorem shows that if one factor has a perfect orthogonal colouring and the other has a square number of vertices, then their tensor graph has a perfect orthogonal colouring.

Theorem 2.2. *If G has n^2 vertices, H has m^2 vertices, and $O\chi(G) = n$, then $O\chi(G \times H) = nm$.*

Proof. Label $V(G) = \{v_k : 0 \leq k < n^2\}$ and $V(H) = \{(u_i, u_j) : 0 \leq i, j < m\}$. Let $f = (f_1, f_2)$ be a proper orthogonal colouring of G where f_1 and f_2 use the colours $\{0, 1, \dots, n-1\}$. It is shown that $g = (g_1, g_2)$ is an orthogonal colouring of $G \times H$ using nm colours, where:

$$g_1((v_k, (u_i, u_j))) = f_1(v_k) + in$$

and

$$g_2((v_k, (u_i, u_j))) = f_2(v_k) + jn.$$

First, it is shown that g has no orthogonal conflicts. Let $v_{k_1}, v_{k_2} \in V(G)$ and let $(u_{i_1}, u_{j_1}), (u_{i_2}, u_{j_2}) \in V(H)$. If $g((v_{k_1}, (u_{i_1}, u_{j_1}))) = g((v_{k_2}, (u_{i_2}, u_{j_2})))$, then:

$$(2.1) \quad f_1(v_{k_1}) + i_1n = f_1(v_{k_2}) + i_2n$$

and

$$(2.2) \quad f_2(v_{k_1}) + j_1n = f_2(v_{k_2}) + j_2n.$$

Without loss of generality, suppose that $i_1 < i_2$. Then it follows that:

$$\begin{aligned} f_1(v_{k_1}) + i_1n &< n + i_1n \\ &\leq i_2n \\ &\leq f_1(v_{k_2}) + i_2n. \end{aligned}$$

Therefore, $f_1(v_{k_1}) + i_1n < f_1(v_{k_2}) + i_2n$, which contradicts equation (2.1), thus $i_1 = i_2$. A similar argument shows that $j_1 = j_2$. Substituting $i_1 = i_2$ and $j_1 = j_2$ into equations (2.1) and (2.2), gives $f_1(v_{k_1}) = f_1(v_{k_2})$ and $f_2(v_{k_1}) = f_2(v_{k_2})$. Hence, $v_{k_1} = v_{k_2}$ because f is an orthogonal colouring of G . Thus, $(v_{k_1}, (u_{i_1}, u_{j_1})) = (v_{k_2}, (u_{i_2}, u_{j_2}))$.

It remains to show that g_1 and g_2 are proper colourings of $G \times H$. Suppose that $v_{k_1}v_{k_2} \in E(G)$ and $(u_{i_1}, u_{j_1})(u_{i_2}, u_{j_2}) \in E(H)$. If $i_1 = i_2 = i$, then since f_1 is a proper colouring of G , $g_1((v_{k_1}, (u_{i_1}, u_{j_1}))) = f_1(v_{k_1}) + in \neq f_1(v_{k_2}) + in = g_1((v_{k_2}, (u_{i_2}, u_{j_2})))$. Thus, there are no colour conflicts between these vertices.

Now, without loss of generality, suppose that $i_1 < i_2$. Then it follows that $g_1((v_{k_1}, (u_{i_1}, u_{j_1}))) = f_1(v_{k_1}) + i_1n < n + i_1n \leq i_2n \leq f_1(v_{k_2}) + i_2n = g_1((v_{k_2}, (u_{i_2}, u_{j_2})))$. Hence, $g_1((v_{k_1}, (u_{i_1}, u_{j_1}))) < g_1((v_{k_2}, (u_{i_2}, u_{j_2})))$. Thus, there are no colour conflicts between these vertices. Therefore, g_1 is a proper colouring. A similar argument shows that g_2 is proper. Thus, g is an orthogonal colouring of $G \times H$. Since $G \times H$ has n^2m^2 vertices and g uses nm colours, $O\chi(G \times H) = nm$. $\square$

Theorem 2.2 provides a method for constructing perfect orthogonal colourings out of graphs that have perfect orthogonal colourings. On the other hand, Theorem 2.1 gives that $K_n \times K_n$ is the maximum graph with n as

its orthogonal chromatic number. Combining these two results provides an upper bound on the orthogonal chromatic number of general tensor graphs.

Corollary 2.3. *If $O_\chi(G) = n$ and $O_\chi(H) = m$, then $O_\chi(G \times H) \leq nm$.*

Proof. Since $O_\chi(G) = n$ and $O_\chi(H) = m$, $G \subseteq K_n \times K_n$ and $H \subseteq K_m \times K_m$ by Theorem 2.1. Therefore, $G \times H \subseteq (K_n \times K_n) \times (K_m \times K_m)$. Since $|V(K_n \times K_n)| = n^2$, $|V(K_m \times K_m)| = m^2$, and $O_\chi(K_n \times K_n) = n$, $O_\chi((K_n \times K_n) \times (K_m \times K_m)) = nm$ by Theorem 2.2. Therefore, $(K_n \times K_n) \times (K_m \times K_m) \subseteq K_{nm} \times K_{nm}$ by Theorem 2.1. Thus, $G \times H \subseteq K_{nm} \times K_{nm}$, and Theorem 2.1 gives that $O_\chi(G \times H) \leq nm$. $\square$

Corollary 2.3 gives an upper bound on the orthogonal chromatic number of tensor graphs in the case where the orthogonal chromatic numbers of the factors are known. However, this upper bound can be far from the exact orthogonal chromatic number. For instance, $O_\chi(K_n) = n$, so Corollary 2.3 gives $O_\chi(K_n \times K_n) \leq n^2$. However, by Theorem 2.1, $O_\chi(K_n \times K_n) = n$.

On the other hand, Corollary 2.3 gives a good upper bound for the bipartite double cover graphs, $G \times K_2$, where G has an optimal orthogonal colouring. These graphs are of interest for other types of colourings [8]. For example, consider the cycle graph C_9 which by [7] has $O_\chi(C_9) = 3$. Then by Corollary 2.3, $O_\chi(C_9 \times K_2) \leq 6$. Which is only one off of the correct orthogonal chromatic number, illustrated in Figure 2.

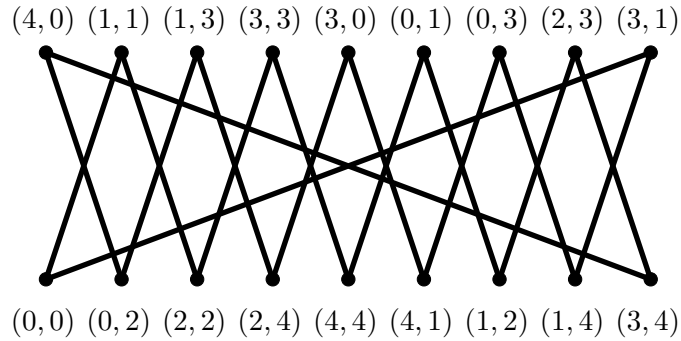


FIGURE 2. Orthogonal colouring of $C_9 \times K_2$

This concludes this section's study of perfect orthogonal colourings of tensor graphs. It remains an open problem to determine the correct orthogonal chromatic number of bipartite double cover graphs. Studying the maximal case, $K_n \times K_2$, may yield an improved upper bound.

3. PERFECT k -ORTHOGONAL COLOURINGS

To start this section, it is shown that if two graphs have a perfect k -orthogonal colouring, then so does their tensor graph. The main idea behind the proof of this result is to create colour classes for the tensor graph out of

the colour classes of the factors. Then, the goal is to show that these new colour classes only share at most one element, and thus give an orthogonal colouring.

Theorem 3.1. *If G has n^2 vertices with $O\chi_k(G) = n$ and H has m^2 vertices with $O\chi_k(H) = m$, then $O\chi_k(G \times H) = nm$.*

Proof. For $0 \leq r < k$ and $0 \leq i < n$, let $G_{r,i}$ be the i th colour class in the r th colouring of G . Then, for $0 \leq r < k$ and $0 \leq j < m$, let $H_{r,j}$ be the j th colour class in the r th colouring of H . Next, let $I_{r,i,j} = \{(u, v) \mid u \in G_{r,i}, v \in H_{r,j}\}$. It will be shown that $C_r = \{I_{r,i,j} \mid 0 \leq i < n, 0 \leq j < m\}$ is a partition of $G \times H$ into nm independent sets. That is, C_r is a proper colouring of $G \times H$ using nm colours.

First, it is shown that $I_{r,i,j}$ is an independent set. Let $(u_1, v_1), (u_2, v_2) \in I_{r,i,j}$. Then $u_1, u_2 \in G_{r,i}$ and $v_1, v_2 \in H_{r,j}$. However, $G_{r,i}$ and $H_{r,j}$ are independent sets, thus $u_1 u_2 \notin E(G)$ and $v_1 v_2 \notin E(H)$. Thus, $(u_1, v_1)(u_2, v_2) \notin E(G \times H)$. Now, let $(u, v) \in V(G \times H)$. Since $\{G_{r,i} \mid 0 \leq i < n\}$ is a partition of G , $u \in G_{r,i}$ for some i . Similarly, $v \in H_{r,j}$ for some j . Therefore, $(u, v) \in I_{r,i,j}$.

Now, suppose that $(u, v) \in I_{r,i_1,j_1}$ and $(u, v) \in I_{r,i_2,j_2}$. If $i_1 \neq i_2$ then $u \in G_{r,i_1}$ and $u \in G_{r,i_2}$. However, this contradicts that $\{G_{r,i} \mid 0 \leq i < n\}$ is a colouring of G . Similarly, if $j_1 \neq j_2$, then $v \in H_{r,j_1}$ and $v \in H_{r,j_2}$. However, this contradicts that $\{H_{r,j} \mid 0 \leq j < m\}$ is a colouring of H . Therefore, there is a unique set $I_{r,i,j}$ that contains (u, v) . Thus, C_r is a proper colouring of $G \times H$ using nm colours.

It remains to show that each of the colourings are mutually orthogonal. Consider I_{r_1,i_1,j_1} and I_{r_2,i_2,j_2} where $r_1 \neq r_2$. If $(u, v) \in I_{r_1,i_1,j_1}$ and $(u, v) \in I_{r_2,i_2,j_2}$, then $u \in G_{r_1,i_1}$ and $u \in G_{r_2,i_2}$. However, $|G_{r_1,i_1} \cap G_{r_2,i_2}| = 1$, so let u be this unique vertex. Similarly, $v \in H_{r_1,j_1}$ and $v \in H_{r_2,j_2}$. However, $|H_{r_1,j_1} \cap H_{r_2,j_2}| = 1$, so let v be this unique vertex. Therefore, there is a unique vertex (u, v) in both I_{r_1,i_1,j_1} and I_{r_2,i_2,j_2} . Hence, each of the C_r are mutually orthogonal. $\square$

Interestingly, the orthogonal colouring created in Theorem 3.1 works for Cartesian graphs as well. The Cartesian graph product of two graphs G and H , denoted by $G \square H$, has vertex set $V(G) \times V(H)$, and two vertices (u_1, v_1) and (u_2, v_2) in $G \square H$ are adjacent if and only if $u_1 = u_2$ and $v_1 v_2 \in E(H)$ or if $v_1 = v_2$ and $u_1 u_2 \in E(G)$. It is also a perfect k -orthogonal colouring for the strong product graph. The strong product of two graphs G and H , denoted $G \boxtimes H$, vertex set $V(G) \times V(H)$ and edge set $E(G \boxtimes H) = E(G \square H) \cup E(G \times H)$.

Corollary 3.2. *If G has n^2 vertices with $O\chi_k(G) = n$ and H has m^2 vertices with $O\chi_k(H) = m$, then $O\chi_k(G \square H) = nm$ and $O\chi_k(G \boxtimes H) = nm$.*

Proof. Let $I_{r,i,j}$ be the same set as in Theorem 3.1. Then, note that $I_{r,i,j}$ is an independent set in $G \square H$ and $G \boxtimes H$. Therefore, this result follows by applying the proof of Theorem 3.1. $\square$

Theorem 3.1 gives a method to construct perfect k -orthogonal colourings when both factors have a perfect k -orthogonal colouring. This is now used to find an upper bound on the k -orthogonal chromatic number of tensor product graphs. Recall Theorem 2.1, which gives a way to reformulate the problem as a subgraph question. Unlike perfect orthogonal colourings, for perfect k -orthogonal colourings, there are multiple graphs required to reformulate the problem.

Caro and Yuster [4] showed that a graph has a perfect k -orthogonal colouring if and only if it is a subgraph of a graph obtained by removing k edge-disjoint K_n -covers from K_{n^2} . Let $K_{n^2}[k]$ denote this family of graphs. Thus, for $k = 2$, Theorem 2.1 gives that $K_{n^2}[2] = K_n \times K_n$. Therefore, using the same argumentation as Corollary 2.3, but using this family of graphs, the following upper bound is obtained.

Corollary 3.3. *If $O_{\chi_k}(G) = n$ and $O_{\chi_k}(H) = m$. then $O_{\chi_k}(G \times H) \leq nm$.*

Proof. Suppose that $O_{\chi_k}(G) = n$. Then $G \subseteq \bar{G}$ and $H \subseteq \bar{H}$ for some $\bar{G} \in K_{n^2}[k]$ and $\bar{H} \in K_{m^2}[k]$. Then, since $\bar{G}$ has n^2 vertices with $O_{\chi_k}(\bar{G}) = n$ and $\bar{H}$ has m^2 vertices with $O_{\chi_k}(\bar{H}) = m$, $O_{\chi_k}(\bar{G} \times \bar{H}) = nm$ by Theorem 3.1. Therefore, since $G \times H \subseteq \bar{G} \times \bar{H}$, $O_{\chi_k}(G \times H) \leq nm$ by restricting the k -orthogonal colouring. $\square$

Corollary 3.3 gives an upper bound on the k -orthogonal chromatic number of tensor graphs in the case where the orthogonal chromatic number of the factors are known. Similar to Corollary 2.3, this gives better upper bounds the closer the k -orthogonal chromatic number of the factors are to being perfect k -orthogonal chromatic numbers. To conclude this paper, the following theorem gives one more method to construct perfect k -orthogonal colourings of tensor graphs.

Theorem 3.4. *If G has n^2 vertices, H has p^2 vertices where p is a prime, and $O_{\chi_k}(G) = n$ with $k \leq p$, then $O_{\chi_k}(G \times H) = np$.*

Proof. Label $V(H) = \{(u_i, u_j) : 0 \leq i, j < p\}$. For $0 \leq r < k$ and $0 \leq s < n$, let $I_{r,s}$ be the s th colour class in the r th colouring of G . Then, for $0 \leq j < p$, let $\bar{I}_{r,s,j} = \{(v, (u_i, u_{(ir+j)(\text{mod } p)})) \mid v \in I_{r,s}, 0 \leq i < p\}$. The goal is to show that $C_r = \{\bar{I}_{r,s,j} \mid 0 \leq s < n, 0 \leq j < p\}$ is a partition of $G \times H$ into np independent sets. That is, C_r is a proper colouring of $G \times H$ using np colours.

First, it is shown that each $\bar{I}_{r,s,j}$ is an independent set. Since each $I_{r,s}$ is an independent set in G , for each $v_1, v_2 \in I_{r,s}$, $v_1 v_2 \notin E(G)$. Thus by the definition of the tensor graph, $(v_1, (u_i, u_{(ir+j)(\text{mod } p)}))(v_2, (u_i, u_{(ir+j)(\text{mod } p)})) \notin E(G \times H)$. Therefore, each $\bar{I}_{r,s,j}$ is an independent set. Next, it is shown that C_r is a partition of $G \times H$.

Consider a vertex $(v, (u_x, u_y))$ in $G \times H$. Since $\{I_{r,s} \mid 0 \leq s < n\}$ is a partition of G , $v \in I_{r,s}$ for some s . Now, notice that for $0 \leq j < p$, $\{(u_i, u_{(ir+j)(\text{mod } p)}) \mid 0 \leq i < p\}$ is a partition of H . Therefore, (u_x, u_y) is

in one of these sets. In particular, this occurs for $i = x$ and $j = y - r$. Therefore, there is a unique set $\bar{I}_{r,s,j}$ that contains $(v, (u_x, u_y))$. Thus, C_r is a partition of $G \times H$ into independent sets.

Now it remains to show that each of the colourings are mutually orthogonal. That is, it remains to show for $r_1 \neq r_2$, s_1, s_2 and j_1, j_2 fixed, that $|\bar{I}_{r_1,s_1,j_1} \cap \bar{I}_{r_2,s_2,j_2}| = 1$. Since $|I_{r_1,s_1} \cap I_{r_2,s_2}| = 1$, let v be this vertex. Therefore, if it can be shown that $|\{(u_{i_1}, u_{i_1 r_1 + j_1}) | 0 \leq i_1 < p\} \cap \{(u_{i_2}, u_{i_2 r_2 + j_2}) | 0 \leq i_2 < p\}| = 1$, then we are done. Since $r_1 \neq r_2$, the only way $(u_{i_1}, u_{i_1 r_1 + j_1}) = (u_{i_2}, u_{i_2 r_2 + j_2})$ is if $i_1 = i_2$. Thus, call this i . Now, consider the following equation.

$$(ir_1 + j_1)(\text{mod } p) = (ir_2 + j_2)(\text{mod } p)$$

This equation simplifies to the following equation.

$$i(r_1 - r_2)(\text{mod } p) = (j_2 - j_1)(\text{mod } p)$$

Since $k \leq p$, and $r_1 \neq r_2$, $r_1 - r_2 \neq 0(\text{mod } p)$. Thus, $r_1 - r_2 = r$ and $j_2 - j_1 = j$. Since p is a prime, $\mathbb{Z}_p$ has no zero divisors. Therefore, $ir(\text{mod } p) = j(\text{mod } p)$ has a unique solution, call this unique solution (i, j) . Thus, $(v, (i, j))$ is the unique element in the $\bar{I}_{r_1,s_1,j_1} \cap \bar{I}_{r_2,s_2,j_2}$. Hence, the colourings are all mutually orthogonal. Since each of these colourings using np colours, and $G \times H$ has $n^2 p^2$ vertices, $O\chi(G \times H) = np$ as desired. $\square$

ACKNOWLEDGEMENT

The author would like to thank his supervisor, Jeannette Janssen, for her helpful comments and suggestions on this paper.

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