



RESOLVABILITY IN HYPERGRAPHS

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ABSTRACT. This article presents an extension of the study of metric and partition dimension to hypergraphs. We give sharp lower bounds for the metric and partition dimension of hypergraphs in general and give exact values under specified conditions.

1. INTRODUCTION

A *hypergraph* H is a pair (V, \mathcal{H}) , where V is a finite non-empty set of vertices and \mathcal{H} is a finite family of distinct non-empty subsets of V , called hyperedges, with

$$\bigcup_{E \in \mathcal{H}} E = V.$$

The “order” and the “size” of H is denoted by n and m , respectively. A hypergraph $K = (V_1, \mathcal{K})$ is a *subhypergraph* of H if and only if $V_1 \subseteq V$ and $\mathcal{K} \subseteq \mathcal{H}$. A hypergraph H is *linear* if for distinct hyperedges $E_i, E_j \in \mathcal{H}$, $|E_i \cap E_j| \leq 1$, so for a linear hypergraph there are no repeated hyperedges of cardinality greater than one. A hypergraph H such that no hyperedge is a subset of any other is called *Sperner*.

A vertex $v \in V$ is *incident* with a hyperedge E of H if $v \in E$. If v is incident with exactly d hyperedges, then we say that the *degree* of v is d ; if all the vertices $v \in V$ has degree r , then H is *r -regular*. Similarly, if there are exactly k vertices incident with a hyperedge E , then we say that the *size* of E is k ; if all the hyperedges $E \in \mathcal{H}$ have size k , then H is *k -uniform*. A graph is simply a 2-uniform hypergraph. A hyperedge E of H is called a *pendant hyperedge* if for $E_i, E_j \in \mathcal{H}$, $E \cap E_i \neq \emptyset$ and $E \cap E_j \neq \emptyset$ implies $(E \cap E_i) \cap (E \cap E_j) \neq \emptyset$. A *path* of length l from a

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vertex v to another vertex u in a hypergraph is a finite sequence of the form $v, E_1, w_1, E_2, w_2, \dots, E_{l-1}, w_{l-1}, E_l, u$ such that $v \in E_1$, $w_i \in E_i \cap E_{i+1}$ for $i = 1, 2, \dots, l-1$ and $u \in E_l$. A hypergraph H is said to be *connected* if there is a path between any two vertices of H . All hypergraphs considered in this paper are connected Sperner hypergraphs.

A hypergraph H is said to be a *hyperstar* if there exists a subset C of vertices such that $E_i \cap E_j = C \neq \emptyset$, for any $E_i, E_j \in \mathcal{H}$. Then C is called the center of the hyperstar. If there exists a sequence of hyperedges E_1, E_2, \dots, E_m in a hypergraph H , then H is said to be (1) a *hyperpath* if $E_i \cap E_j \neq \emptyset$ if and only if $|i - j| = 1$; (2) a *hypercycle* if $E_i \cap E_j \neq \emptyset$ if and only if $i - j \in \{1, -1\} \pmod{m}$ for $m \geq 3$. A connected hypergraph H with no hypercycle is called a *hypertree*. A subhypertree of a hypertree H with edge set, say $\mathcal{E} = \{E_{p_1}, E_{p_2}, \dots, E_{p_l}\} \subset \mathcal{H}$, is called a *branch* of H if E_{p_1} (say) is the only hyperedge such that, for $E_i, E_j \in \mathcal{H} - \mathcal{E}$, $E_{p_1} \cap E_i \neq \emptyset$ and $E_{p_1} \cap E_j \neq \emptyset$ implies $(E_{p_1} \cap E_i) \cap (E_{p_1} \cap E_j) \neq \emptyset$ as well as $E_{p_i} \cap E_j = \emptyset$ for each $2 \leq i \leq l$ and for all $E_j \in \mathcal{H} - \mathcal{E}$. The hyperedge E_{p_1} is then called the *joint* of the branch.

An ordered set W of vertices of a connected graph G is called a resolving set for G if for every two distinct vertices $u, v \in V(G)$, there is a vertex $w \in W$ such that $d(u, w) \neq d(v, w)$. A resolving set of minimum cardinality is called a basis for G and the number of vertices in a basis is called the metric dimension of G , denoted by $\dim(G)$. An ordered t -partition $\Pi = \{S_1, S_2, \dots, S_t\}$ of $V(G)$ is called a resolving partition if for every two distinct vertices $u, v \in V(G)$, there is a set S_i in Π such that $d(u, S_i) \neq d(v, S_i)$, where

$$d(v, s) = \min_{s \in S} d(u, s).$$

The minimum t for which there is a resolving t -partition of $V(G)$ is called the partition dimension of G , denoted by $\text{pd}(G)$. In this article, we consider hypergraphs in the context of metric dimension and partition dimension, which are defined in Sections 2 and 3, respectively. We give sharp lower bounds for the metric and partition dimension of graphs. The metric dimension of some well-known families of hypergraphs such as hyperpaths, hypertrees and k -uniform linear hypercycles is investigated. Further, we find the metric and partition dimensions of 3-uniform linear hypercycles as well as the partition dimension of k -uniform hyperpath.

2. METRIC DIMENSION OF HYPERGRAPHS

The metric dimension of a graph was studied by Slater [14] and independently by Harary and Melter [6]. It is a parameter that has appeared in

various applications, as diverse as combinatorial optimization, pharmaceutical chemistry, robot navigation and sonar. In recent years, considerable literature has been developed (see [12, 4, 10, 1, 7, 8, 13, 11]). The problem of determining whether $\dim(H) < M$ ($M > 0$), where H is a simple graph, is an NP-complete problem [5, 13]. The metric dimension of a hypergraph H is defined as follows:

The *distance* between any two vertices v and u of a hypergraph H , $d(v, u)$, is the length of the shortest path between them and $d(v, u) = 0$ if and only if $v = u$. The *diameter* of H is the maximum distance between the vertices of H , and is denoted by $\text{diam}(H)$. Two vertices u and v of H are said to be “diametral” vertices if $d(u, v) = \text{diam}(H)$. The *representation*, $r(v|W)$, of a vertex v of H with respect to an ordered set $W = \{w_1, w_2, \dots, w_q\} \subseteq V$ is the q -tuple $r(v|W) = (d(v, w_1), d(v, w_2), \dots, d(v, w_q))$. The set W is called a *resolving set* for a hypergraph H if $r(v|W) \neq r(u|W)$ for any two different vertices $v, u \in V$. A resolving set with minimum cardinality is called a *basis* for H and that minimum cardinality is called the *metric dimension* of H , denoted by $\dim(H)$.

To determine whether a given set $W \subseteq V$ is a resolving set for a hypergraph H , W needs only to be verified for the vertices in $V - W$ since every vertex $w \in W$ is the only vertex of H whose distance from w is 0.

For a hypergraph H , let the set of hyperedges in H be $\mathcal{H} = \{E_1, E_2, \dots, E_m\}$. Let us denote the set of all the vertices in the set $E_{i_1} \cap E_{i_2} \cap \dots \cap E_{i_d}$ having degree d by $\mathcal{C}(i_1, i_2, \dots, i_d)$, where each $i_j \in \{1, 2, \dots, m\}$ is distinct. That is

$$\mathcal{C}(i_1, i_2, \dots, i_d) = \{v \in V \mid v \in E_{i_1} \cap E_{i_2} \cap \dots \cap E_{i_d} \text{ and } \deg(v) = d\}.$$

Then, the collection of all such classes defines a partition of the vertex set V . Let $\eta(i_1, i_2, \dots, i_d) = |\mathcal{C}(i_1, i_2, \dots, i_d)| - 1$ whenever $\mathcal{C}(i_1, i_2, \dots, i_d) \neq \emptyset$, otherwise we take $\eta(i_1, i_2, \dots, i_d) = 0$.

Example. The hypergraph H shown in the Figure 1 with set of vertices $V = \{v_1, v_2, v_3, \dots, v_9\}$ and hyperedges $E_1 = \{v_1, v_2, v_3, v_4\}$, $E_2 = \{v_3, v_4, v_5, v_6, v_7\}$, $E_3 = \{v_6, v_7, v_8, v_9\}$, $E_4 = \{v_2, v_3, v_9\}$.

In the graph H , $\mathcal{C}(1) = \{v_1\}$, $\mathcal{C}(2) = \{v_5\}$, $\mathcal{C}(3) = \{v_8\}$, $\mathcal{C}(4) = \emptyset$; $\mathcal{C}(1, 2) = \{v_4\}$, $\mathcal{C}(1, 3) = \emptyset$, $\mathcal{C}(1, 4) = \{v_2\}$, $\mathcal{C}(2, 3) = \{v_6, v_7\}$, $\mathcal{C}(2, 4) = \emptyset$, $\mathcal{C}(3, 4) = \{v_9\}$; $\mathcal{C}(1, 2, 3) = \emptyset$, $\mathcal{C}(1, 2, 4) = \{v_3\}$, $\mathcal{C}(1, 3, 4) = \emptyset$, $\mathcal{C}(2, 3, 4) = \emptyset$ and $\mathcal{C}(1, 2, 3, 4) = \emptyset$. Note that the set of all given classes form a partition for V .

Thus, we have the following straightforward proposition:

Proposition. For any two distinct vertices $u, v \in \mathcal{C}(i_1, i_2, \dots, i_d)$, we have $d(u, w) = d(v, w)$ for any $w \in V - \{u, v\}$.

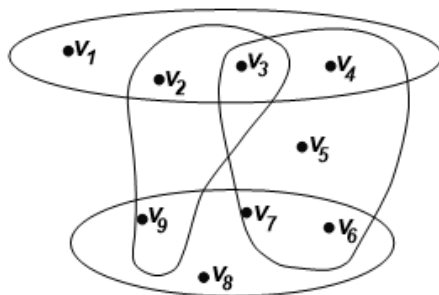


FIGURE 1. Hypergraph H

Thus, we extract the following Lemma related to the resolving set for H :

Lemma. *If $u, v \in \mathcal{C}(i_1, i_2, \dots, i_d)$ and $W \subseteq V$ resolves H , then at least one of the vertices u and v is in W . Moreover, if $u \in W$ and $v \notin W$, then $(W - \{u\}) \cup \{v\}$ also resolves H .*

Now, we can establish a lower bound for the metric dimension of a hypergraph in the following result.

Proposition. *For any hypergraph H with m hyperedges,*

$$\dim(H) \geq \sum_{j=1}^m \sum_{i_1 < \dots < i_j} \eta(i_1, i_2, \dots, i_j).$$

Proof. It follows from the fact that if there are $|\mathcal{C}(i_1, i_2, \dots, i_d)|$ number of vertices of degree d in $E_{i_1} \cap E_{i_2} \cap \dots \cap E_{i_d}$, then Lemma 2.3 yields that at least $\eta(i_1, i_2, \dots, i_d)$ vertices should belong to any basis W . \square

Remark. *By Proposition 2.4, it is clear that, in order to obtain a basis of any hypergraph H , it suffices to consider only one vertex, say v_{i_1, i_2, \dots, i_d} , from each class $\mathcal{C}(i_1, i_2, \dots, i_d)$ if $\mathcal{C}(i_1, i_2, \dots, i_d) \neq \emptyset$. We call this vertex, a representative vertex of the class $\mathcal{C}(i_1, i_2, \dots, i_d)$. We denote the set of all representative vertices in a hypergraph H by $R(H)$, and hence we always have, $V - R(H) \subseteq W$ for any basis W of H .*

Now we discuss some classes of hypergraphs for which the equality holds in the Proposition 2.4.

Theorem. *For any hypergraph H with m hyperedges, if $\eta(i) \neq 0$ for all $E_i \in E(H)$, then*

$$\dim(H) = \sum_{j=1}^m \sum_{i_1 < \dots < i_j} \eta(i_1, i_2, \dots, i_j).$$

Moreover, there are

$$\prod_{j=1}^m \prod_{i_1 < \dots < i_j} (\eta(i_1, i_2, \dots, i_j) + 1)$$

basis for H .

Proof. Consider $W = V - R(H)$, we have to show that W is a basis for H . Take any two different vertices $v, v' \in R(H)$. Since both the vertices v and v' are representative vertices of different classes, there exists a hyperedge E_j such that $v' \in E_j$ and $v \notin E_j$. It follows from $\eta(i) \neq 0$ that there exists a vertex of degree one $w_j \in V$ such that $w_j \in E_j \cap W$. Clearly, $d(v', w_j) = 1$ and $d(v, w_j) \neq 1$, hence W is a basis for H . Further, by Lemma 2.3, there are $\prod_{j=1}^m \prod_{i_1 < \dots < i_j} (\eta(i_1, i_2, \dots, i_j) + 1)$ such W . \square

If H is a linear hypergraph with m hyperedges, then $\eta(i, j) = 0$ for every $E_i, E_j \in \mathcal{H}$. Thus, we have the following corollary:

Corollary. *Let H be a linear hypergraph with m hyperedges and $\eta(i) \neq 0$ for all $E_i \in \mathcal{H}$. Then $\dim(H) = \sum_{i=1}^m \eta(i)$.*

We give two examples which show that the condition in Theorem 2.6 cannot be relaxed in general.

Example. *Let H be a hypergraph with vertex set $V = \{v_1, v_2, v_3, v_4\}$ and edge set $\mathcal{H} = \{E_1, E_2\}$, where $E_1 = \{v_1, v_2, v_3\}$ and $E_2 = \{v_3, v_4\}$. Clearly, $\eta(2) = 0$ so H does not satisfy the condition of Theorem 2.6. Without loss of generality, we can take the set of representative vertices $R(H) = \{v_1, v_3, v_4\}$, and hence $W = V - R(H) = \{v_2\}$. But, W is not a resolving set for H since $r(v_1|W) = r(v_3|W)$. In fact, $\dim(H) = 2 > 1$.*

Example. *Let H be a hypergraph with vertex set $V = \{v_1, v_2, v_3, v_4, v_5, v_6\}$ and edge set $\mathcal{H} = \{E_1, E_2, E_3\}$, where $E_1 = \{v_1, v_2, v_3, v_4\}$, $E_2 = \{v_3, v_4, v_5, v_6\}$ and $E_3 = \{v_1, v_2, v_5, v_6\}$. Clearly, $\eta(i) = 0$ for all $i = 1, 2, 3$ and $\eta(1, 2) = \eta(2, 3) = \eta(3, 1) \neq 0$. Without loss of generality, we can take the set of representative vertices $R(H) = \{v_1, v_3, v_5\}$, and hence $W = V - R(H) = \{v_2, v_4, v_6\}$. But, W is not a resolving set for H since $r(v_1|W) = r(v_3|W) = r(v_5|W)$. In fact, $\dim(H) = 5 > 3$.*

However, the condition in Theorem 2.6 can be reduced in some special cases as shown in the following result.

Theorem. *Let H be a hyperpath with m hyperedges E_1, E_2, \dots, E_m in a canonical way. Then*

$$\dim(H) = \sum_{i=1}^m \eta(i) + \sum_{i=1}^{m-1} \eta(i, i+1)$$

if both $\eta(1)$ and $\eta(m)$ are non-zero.

Proof. Let $W = V - R(H)$. Then it follows from the facts $\eta(1) \neq 0$ and $\eta(m) \neq 0$ that there exists a vertex of degree one $w_1 \in E_1 \cap W$ and there exists a vertex of degree one $w_m \in E_m \cap W$. To prove the theorem, we only have to show that the representative vertices are resolved by the set W , and it yields from the fact that for any $1 \leq j \leq m$, we have $(d(v_j, w_1), d(v_j, w_m)) = (j, m - j + 1)$, and for any $1 \leq j < m - 1$, we have $(d(v_{j,j+1}, w_1), d(v_{j,j+1}, w_m)) = (j, m - j)$. \square

Theorem. *Let H be a hypertree with m hyperedges and let $E_{p_1}, E_{p_2}, \dots, E_{p_t}$ be its pendant hyperedges. Then*

$$\dim(H) = \sum_{j=1}^m \sum_{i_1 < \dots < i_j} \eta(i_1, i_2, \dots, i_j)$$

if $\eta(p_s) \neq 0$ for all $s = 1, 2, \dots, t$.

Proof. Consider $W = V - R(H)$, similarly as in the proof of Theorem 2.6, again we have to show that W is a basis for H . Take any two different vertices $v, v' \in R(H)$, then both vertices are representative of two different classes, and hence there exists a hyperedge E_j such that $v' \in E_j$ but $v \notin E_j$. Now, consider a hyperpath contained in the hypertree H which starts and ends at the pendant hyperedges and contains both v and E_j . By using the proof of Theorem 2.10, it can be seen that the vertices v and v' has different representations for W , which proves the theorem. \square

In a k -uniform ($k \geq 3$) linear hyperstar H is a special case of hypertree in which $\eta(i) \neq 0$ for all $E_i \in \mathcal{H}$, so we have the following corollary:

Corollary. *If H is a k -uniform ($k \geq 3$) linear hyperstar with m ($m \geq 2$) hyperedges, then $\dim(H) = m(k - 2)$.*

Consider a k -uniform ($k \geq 4$) linear hypercycle $\mathcal{C}_{m,k}$ with m hyperedges. Then $\eta(i) \neq 0$ for all edges E_i of $\mathcal{C}_{m,k}$. By Corollary 2.7, $\dim(\mathcal{C}_{m,k}) = m(k - 3)$. A 3-uniform linear hypercycle with four hyperedges is shown in Figure 2. For the case $k = 3$, we have $\eta(i) = 0$ for all $E_i \in \mathcal{H}$, hence the lower bound given in Proposition 2.4 is zero and every vertex in $\mathcal{C}_{m,3}$ is the representative vertex. We discuss this case in the following result:

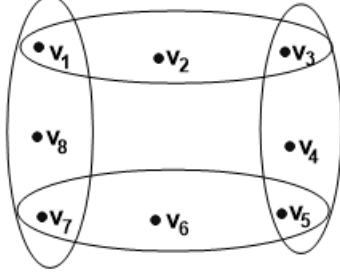


FIGURE 2. A 3-uniform linear hypercycle with 4 hyperedges

Theorem. Let $\mathcal{C}_{m,3}$ be a 3-uniform linear hypercycle with m hyperedges. Then $\dim(\mathcal{C}_{3,3}) = 2$ and for all $m \geq 4$,

$$\dim(\mathcal{C}_{m,3}) = \begin{cases} 2, & \text{if } k \text{ is even,} \\ 3, & \text{if } k \text{ is odd.} \end{cases}$$

Proof. In $\mathcal{C}_{m,3}$, each $v_j \in E_j$ represents a vertex of degree one and $v_{j,j+1} \in E_j \cap E_{j+1}$ with $v_{m,m+1} = v_{m,1}$. Clearly, $\dim(\mathcal{C}_{m,3}) > 1$ for any m .

If m is even, then we take $W = \{v_1, v_{\frac{m}{2}}\}$.

For $1 < j < \frac{m}{2}$, we have $r(v_j|W) = (j, \frac{m}{2} - j + 1)$ and for $1 \leq j < \frac{m}{2}$, $r(v_{j,j+1}|W) = (j, \frac{m}{2} - j)$. Now, if $\frac{m}{2} + 1 \leq j < k$, then $r(v_j|W) = (m + 2 - j, j - \frac{m}{2} + 1)$ and $r(v_{j,j+1}|W) = (m + 1 - j, j - \frac{m}{2} + 1)$ with $r(v_m|W) = (2, \frac{m}{2} + 1)$, $r(v_{\frac{m}{2}, \frac{m}{2}+1}|W) = (\frac{m}{2}, 1)$ and $r(v_{m,1}|W) = (1, \frac{m}{2})$. It is easy to see that the representations of all the vertices with respect to W are distinct, hence W forms a basis for $\mathcal{C}_{m,3}$ and $\dim(\mathcal{C}_{m,3}) = 2$.

For the special case when $m = 3$, the set $W = \{v_1, v_2\}$ forms a basis for $\mathcal{C}_{3,3}$. Hence $\dim(\mathcal{C}_{3,3}) = 2$.

If $m > 3$ is odd, then we first show that $\dim(\mathcal{C}_{m,3}) > 2$. Suppose to the contrary that $\dim(\mathcal{C}_{m,3}) = 2$ and let W be a basis of $\mathcal{C}_{m,3}$. Let us call the vertices $v_{i,i+1}$, $i \in \{1, 2, \dots, m\}$, of $\mathcal{C}_{m,3}$, the common vertices. We have the following three possibilities:

- (1) W contains two common vertices. Without loss of generality, we may assume that one vertex is $v_{1,2}$ and the second vertex is $v_{j,j+1}$ ($2 \leq j \leq m$). Then $r(v_{j+1}|W) = r(v_{j+1,j+2}|W)$, for $2 \leq j < \frac{m+1}{2}$; $r(v_2|W) = r(v_{m,1}|W)$, for $j = \frac{m+1}{2}$; $r(v_1|W) = r(v_{2,3}|W)$, for $j = \frac{m+1}{2} + 1$ and $r(v_j|W) = r(v_{j-1,j}|W)$, for $\frac{m+1}{2} + 1 < j \leq m$, a contradiction.
- (2) W contains one common vertex. Without loss of generality, we may assume that one vertex is $v_{1,2}$ and the second vertex is

v_j ($1 \leq j \leq m$). Then $r(v_{j+1}|W) = r(v_{j+1,j+2}|W)$, for $1 \leq j < \frac{m+1}{2}$; $r(v_1|W) = r(v_{m,1}|W)$, for $j = \frac{m+1}{2}$; $r(v_1|W) = r(v_2|W)$, for $j = \frac{m+1}{2} + 1$ and $r(v_2|W) = r(v_{2,3}|W)$, for $\frac{m+1}{2} + 1 < j \leq m$, a contradiction.

- (3) W contains no common vertex. Without loss of generality, we may assume that one vertex is v_1 and the second vertex is v_j ($2 \leq j \leq m$). Then it will lead to a contradiction as in (1).

Now, we will show that $\dim(\mathcal{C}_{m,3}) \leq 3$. Take $W = \{v_1, v_2, v_{\frac{m+1}{2}}\}$. We note that, $r(v_{1,2}|W) = (1, 1, \frac{m-1}{2})$ and

$$r(v_j|W) = \begin{cases} (j, j-1, \frac{m+1}{2} - j + 1), & \text{for } 2 < j < \frac{m+1}{2}, \\ (\frac{m+1}{2}, \frac{m+1}{2}, 2), & \text{for } j = \frac{m+1}{2} + 1, \\ (m-j+2, m-j+3, j - \frac{m-1}{2}), & \text{for } \frac{m+1}{2} + 1 < j \leq k, \end{cases}$$

$$r(v_{j,j+1}|W) = \begin{cases} (j, j-1, \frac{m+1}{2} - i), & \text{for } 2 \leq j < \frac{m+1}{2}, \\ (\frac{m+1}{2}, \frac{m-1}{2}, 1), & \text{for } j = \frac{m+1}{2}, \\ (m-j+1, m-j+2, j - \frac{m-1}{2}), & \text{for } \frac{m+1}{2} < j \leq k. \end{cases}$$

One can see that all the vertices of $V(\mathcal{C}_{m,3}) - W$ have distinct representations. This implies that $\dim(\mathcal{C}_{m,3}) = 3$ when $m > 3$ is odd. \square

The primal graph, $\text{prim}(H)$, of a hypergraph H is the graph with vertex set V such that vertices x and y of $\text{prim}(H)$ are adjacent if and only if x and y are contained in the same hyperedge. A loop on a vertex in $\text{prim}(H)$ will exist if it is the only vertex incident with a hyperedge. The *middle graph*, $M(H)$, of H is a subgraph of $\text{prim}(H)$ obtained by deleting loops and parallel edges. Since the adjacencies between the vertices in $\text{prim}(H)$ are due to the adjacencies in the hypergraph H , so determining the length of a path between two vertices u and v in $\text{prim}(H)$ is equivalent to determine the length of a path between the vertices u and v in H . This fact yields the following result:

Theorem. *Let H be a hypergraph. Then*

$$\dim(H) = \dim(\text{prim}(H)) = \dim(M(H)).$$

The *dual* of $H = (\{v_1, v_2, \dots, v_n\}, \{E_1, E_2, \dots, E_m\})$, denoted by H^* , is the hypergraph whose vertices are $\{E_1, E_2, \dots, E_m\}$ corresponding to the hyperedges of H and with hyperedges $V_i = \{E_j : v_i \in E_j \text{ in } H\}$, where $i = 1, 2, \dots, n$. In other words, the dual H^* swaps the vertices and hyperedges of H . The primal graph of the dual H^* of a hypergraph H is not a simple graph, in this case, the middle graph of H^* is a simple graph. Since the dual H^* of a hypergraph H is also a hypergraph so $\dim(H^*) = \dim(M(H^*))$. Moreover, the middle graph of dual H^* of H is (1) a simple path P_n if

and only if H is a hyperpath; (2) a simple cycle C_n if and only if H is a hypercycle. In [4], all the simple connected graphs having metric dimension one were characterized by proving the result “ $\dim(G)$ is one if and only if G is a simple path P_n ($n \geq 1$)”. It is straightforward that the metric dimension of the dual of a hypergraph H^* is 1 if and only if H is a hyperpath. In [6], it was shown that the metric dimension of a simple cycle C_n ($n \geq 3$) is 2. So, the metric dimension of the dual H^* of a hypercycle is 2.

3. PARTITION DIMENSION OF HYPERGRAPHS

Possibly to gain insight into the metric dimension, Chartrand et al. introduced the notion of a resolving partition and partition dimension [3, 2]. To define the partition dimension, the distance $d(v, S)$ between a vertex v in H and $S \subseteq V$ is defined as

$$\min_{s \in S} d(v, s).$$

Let $\Pi = \{S_1, S_2, \dots, S_t\}$ be an ordered t -partition of V and v be any vertex of H . Then the *representation*, $r(v|\Pi)$, of v with respect Π is the t -tuple $r(v|\Pi) = (d(v, S_1), d(v, S_2), \dots, d(v, S_t))$. The partition Π is called a *resolving partition* for a hypergraph if $r(v|\Pi) \neq r(u|\Pi)$ for any two distinct vertices $v, u \in V$. The *partition dimension* of a hypergraph H is the cardinality of a minimum resolving partition, denoted by $\text{pd}(H)$.

From the definition of a resolving partition, it can be observed that the property of a given partition Π of a hypergraph H to be a resolving partition of H can be verified by investigating the pairs of vertices in the same class. Indeed, $d(x, S_i) = 0$ for every vertex $x \in S_i$ but $d(x, S_j) \neq 0$ with $j \neq i$. It follows that $x \in S_i$ and $y \in S_j$ are resolved either by S_i or S_j for every $i \neq j$. From Proposition 2.2, we have the following lemma:

Lemma. *Let Π be a resolving partition of V . If $u, v \in \mathcal{C}(i_1, i_2, \dots, i_d)$, then u and v belong to distinct classes of Π .*

The following result gives the lower bound for the partition dimension of hypergraphs.

Proposition. *Let H be a Sperner hypergraph with $m \geq 2$ hyperedges. Then $\text{pd}(H) \geq \lambda + 1$, where $\lambda = \max |\mathcal{C}(i_1, i_2, \dots, i_d)|$ in H .*

Proof. Since $\lambda = \max |\mathcal{C}(i_1, i_2, \dots, i_d)|$ in H , by Lemma 3.1, we have at least λ disjoint classes in any resolving partition $S_1, S_2, \dots, S_\lambda$ of V . Since H is Sperner, there exists an edge E of H such that $\mathcal{C}(i_1, i_2, \dots, i_d) \subset E$. If $u \in E - \mathcal{C}(i_1, i_2, \dots, i_d)$, $v \in \mathcal{C}(i_1, i_2, \dots, i_d)$ and $u, v \in S_i$ for some $1 \leq i \leq \lambda$, then $r(u|\Pi) = r(v|\Pi)$ which is a contradiction. Thus $\text{pd}(H) \geq \lambda + 1$. \square

The lower bound given in Proposition 3.2 is attainable for k -uniform linear hyperpaths as proved in Theorem 3.4.

A 2-uniform hypercycle $\mathcal{C}_{m,2}$ is a simple connected cycle on n vertices and it was shown that the partition dimension of a simple connected cycle is 3 [2], so $\text{pd}(\mathcal{C}_{m,2}) = 3$. In the next result, we investigate the partition dimension of a k -uniform hypercycle $\mathcal{C}_{m,k}$ for $m \geq 3$ and $k \geq 4$.

Theorem. *For $k \geq 4$, let $\mathcal{C}_{m,k}$ be a k -uniform linear hypercycle with $m \geq 3$ hyperedges. Then $\text{pd}(\mathcal{C}_{m,k}) = k$.*

Proof. Firstly, we discuss the case when $k = 3$. In this case, let $v_j \in E_j$ be a vertex of degree one and $v_{j,j+1} \in E_j \cap E_{j+1}$ with $v_{m,m+1} = v_{m,1}$ be a vertex of 2 in $\mathcal{C}_{m,3}$. If we put all the vertices of $\mathcal{C}_{m,3}$ into two classes S_1 and S_2 , then they do not form a resolving partition Π of $V(\mathcal{C}_{m,3})$, because for some hyperedge $E_j \in \mathcal{H}$ such that $E_i \cap S_1 \neq \emptyset$, $E_j \cap S_2 \neq \emptyset$ and $u, v \in S_i$, $r(u|\Pi) = r(v|\Pi)$, which is a contradiction. Thus, $\text{pd}(\mathcal{C}_{m,3}) \geq 3$. On the other hand, $\text{pd}(\mathcal{C}_{m,3}) \leq 3$, because we have a resolving partition of cardinality 3 for $(\mathcal{C}_{m,3})$ in each of the following cases:

For $m \equiv 0 \pmod{6}$, we have a resolving partition for $\text{pd}(\mathcal{C}_{m,3})$ as

$$\Pi = \left\{ \left\{ v_{m,1}, \dots, v_{\frac{1}{3}m, \frac{1}{3}m+1} \right\}, \left\{ v_{\frac{1}{3}m+1}, \dots, v_{\frac{2}{3}m, \frac{2}{3}m+1} \right\}, \left\{ v_{\frac{2}{3}m+1}, \dots, v_m \right\} \right\}.$$

For $m \equiv 1, 4 \pmod{6}$, we have a resolving partition for $\text{pd}(\mathcal{C}_{m,3})$ as

$$\Pi = \left\{ \left\{ v_{m,1}, \dots, v_{\frac{1}{3}(m+2)} \right\}, \left\{ v_{\frac{1}{3}(m+2), \frac{1}{3}(m+2)+1}, \dots, v_{\frac{2}{3}(m+2)-1, \frac{2}{3}(m+2)} \right\}, \right. \\ \left. \left\{ v_{\frac{2}{3}(m+2)}, \dots, v_m \right\} \right\}.$$

For $m \equiv 2 \pmod{6}$, we have a resolving partition for $\text{pd}(\mathcal{C}_{m,3})$ as

$$\Pi = \left\{ \left\{ v_{m,1}, \dots, v_{\frac{1}{3}(m+1)} \right\}, \left\{ v_{\frac{1}{3}(m+1), \frac{1}{3}(m+1)+1}, \dots, v_{\frac{2}{3}(m+1)-1, \frac{2}{3}(m+1)} \right\}, \right. \\ \left. \left\{ v_{\frac{2}{3}(m+1)}, \dots, v_m \right\} \right\}.$$

For $m \equiv 3 \pmod{6}$, we have a resolving partition for $\text{pd}(\mathcal{C}_{m,3})$ as

$$\Pi = \left\{ \left\{ v_1, \dots, v_{\frac{1}{3}m+1} \right\}, \left\{ v_{\frac{1}{3}m+1, \frac{1}{3}m+2}, \dots, v_{\frac{2}{3}m+1} \right\}, \right. \\ \left. \left\{ v_{\frac{2}{3}m+1, \frac{2}{3}m+2}, \dots, v_{m,1} \right\} \right\}.$$

For $m \equiv 5 \pmod{6}$, we have a resolving partition for $\text{pd}(\mathcal{C}_{m,3})$ as

$$\Pi = \left\{ \left\{ v_1, \dots, v_{\frac{1}{3}(m+1), \frac{1}{3}(m+1)+1} \right\}, \left\{ v_{\frac{1}{3}(m+1)+1}, \dots, v_{\frac{2}{3}(m+1)} \right\}, \right. \\ \left. \left\{ v_{\frac{2}{3}(m+1), \frac{2}{3}(m+1)+1}, \dots, v_{m,1} \right\} \right\}.$$

Now, for $k \geq 4$, let us denote the vertices of degree one in j^{th} hyperedge E_j ($1 \leq j \leq m$) by $v_1^j, v_2^j, \dots, v_{k-2}^j$, and let the vertex $v_{j,j+1} \in E_j \cap E_{j+1}$ be

of degree 2 in $\mathcal{C}_{m,k}$. Construct a partition Π with k parts of the vertex set of $\mathcal{C}_{m,k}$ as follows: Copy the resolving partition for $\mathcal{C}_{m,3}$ and for the other vertices that are in the middle of hyperedges, put i^{th} vertex (for $i > 1$) in the $(i + 2)^{\text{th}}$ part. Then it is an easy exercise to verify, by following Lemma 3.1 and Proposition 3.2, that Π is a minimum resolving partition for $\mathcal{C}_{m,3}$, and it completes the proof. \square

Every 2-uniform linear hyperpath is a simple path graph whose partition dimension is 2 as shown in [3]. Now, we generalize this result for more than two uniform linear hyperpaths as follows.

Theorem. *For $k \geq 3$, let H be a k -uniform linear hyperpath with m hyperedges. Then $\text{pd}(H) = k$.*

Proof. Let H be a k -uniform linear hyperpath. Then it is a routine exercise to verify that a partition $\Pi = \{S_1, S_2, \dots, S_k\}$ of V , where each S_i , $1 \leq i \leq k - 1$, contains the i^{th} vertex of every hyperedge of H and S_k contains the k^{th} vertex of the m^{th} hyperedge, is a resolving partition for H . It follows that $\text{pd}(H) \leq k$.

Since H is Sperner and $\max |C(i_1, i_2, \dots, i_d)| = k - 1$ in H , so Proposition 3.2 implies that $\text{pd}(H) \geq k$. \square

The *rank* of a hypergraph H is the maximum number of vertices in a hyperedge. One might think that the partition dimension of H is always greater than or equal to the rank of H . This is true for a k -uniform linear hyperpath and a k -uniform linear hypercycle $\mathcal{C}_{m,3}$. But, in general, it is not true as shown in the following example:

Example. *Let H be a hypergraph with vertex set $V = \{v_i : 1 \leq i \leq 11\}$ and edge set $\mathcal{H} = \{E_1, E_2\}$, where $E_1 = \{v_i; 1 \leq i \leq 7\}$ and $E_2 = \{v_i; 6 \leq i \leq 11\}$. Clearly, $\text{rank}(H) = 7$, $\lambda = 5$ and*

$$\Pi = \{S_i = \{v_i, v_{i+5}\}; 1 \leq i \leq 5, S_6 = \{v_{11}\}\}$$

is a minimum resolving partition of V . This implies that $\text{pd}(H) = 6 \neq \text{rank}(H)$.

Likewise the results on the metric dimension of the primal, we have the following result on the partition dimension of the primal graph of a hypergraph.

Theorem. *Let H be a hypergraph. Then $\text{pd}(H) = \text{pd}(\text{prim}(H))$.*

Let H^* be the dual of a hypergraph H , since H^* is also a hypergraph, so $\text{pd}(H^*) = \text{pd}(M(H^*))$. Since it was shown that the simple paths P_n are the only graphs with $\text{pd}(P_n) = 2$ [3] and the partition dimension of the simple cycles C_n is 3. Therefore, partition dimension of dual H^* of a hyperpath is

2 and partition dimension of dual H^* of a hypercycle is 3.

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