



HAMILTONIAN-CONNECTEDNESS OF TRIANGULATIONS WITH FEW SEPARATING TRIANGLES

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ABSTRACT. We prove that 3-connected plane triangulations containing a single edge contained in all separating triangles are hamiltonian-connected. As a direct corollary we have that 3-connected plane triangulations with at most one separating triangle are hamiltonian-connected. In order to show bounds on the strongest form of this theorem, we proved that for any $s \geq 4$ there are 3-connected triangulation with s separating triangles that are not hamiltonian-connected. We also present computational results which show that all ‘small’ 3-connected triangulations with at most 3 separating triangles are hamiltonian-connected.

1. INTRODUCTION

Plane triangulations – sometimes also called maximal planar graphs – are plane graphs where all faces, including the outer face, are triangles. A hamiltonian cycle, resp. path, is a cycle, resp. path, that visits each vertex exactly once. A graph is hamiltonian if it contains a hamiltonian cycle. A graph is hamiltonian-connected if for every pair of vertices there exists a hamiltonian path having these two vertices as end points. Clearly, being hamiltonian-connected implies being hamiltonian for graphs on $n \geq 3$ vertices.

In 1931, Whitney [9] proved that every 4-connected plane triangulation is hamiltonian. The condition of being 4-connected was later relaxed by Böhme, Harant, and Tkáč who proved that 3-connected triangulations with at most two separating triangles are hamiltonian [1]. An even stronger result of this form is given by Jackson and Yu [4]. They show that a nonhamiltonian triangulation must contain at least 4 separating triangles and even then they have to be in a specific configuration.

A stronger result of a different form was proved by Thomassen who showed that every 4-connected plane graph is hamiltonian-connected [8]. Ozeki and Vrána [6] recently proved an even stronger result in function

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of a property called *k-edge-hamiltonian-connected*. A graph G is *k-edge-hamiltonian-connected* if for any $X \subset \{x_1x_2 : x_1, x_2 \in V(G), x_1 \neq x_2\}$ such that $1 \leq |X| \leq k$ and the graph induced by X on $V(G)$ is a forest in which each component is a path, $G \cup X$ has a hamiltonian cycle containing all edges in X , where $G \cup X$ is the graph obtained from G by adding all edges in X . Note that being 1-edge-hamiltonian-connected is equivalent to being hamiltonian-connected for graphs on $n \geq 3$ vertices, since a hamiltonian cycle in $G \cup \{xy\}$ containing the edge xy corresponds to a hamiltonian path between x and y in G . Ozeki and Vrána proved the following theorem.

Theorem 1.1 ([6]). *Every 4-connected plane graph is 2-edge-hamiltonian-connected.*

The main result of the next section follows from this theorem, but it is actually sufficient to use the following result by Sanders.

Theorem 1.2 ([7]). *For every 4-connected plane graph G with two vertices x, y and an edge e with $e \neq xy$, G contains a hamiltonian path from x to y through e .*

In Section 2, we will show that plane triangulations with only one separating triangle are hamiltonian-connected. We will do this by showing a stronger result about 3-connected triangulations having an edge that is contained in all separating triangles. In Section 3 and Section 4 we investigate how much stronger this theorem could be made using an investigation similar to the one in [2]. We first show that starting with 4 separating triangles there always exist 3-connected triangulations that are not hamiltonian-connected, and we further specify this in terms of the decomposition tree of a 3-connected triangulation. Finally, in Section 4 we give the results of some computer searches to check the hamiltonian-connectedness of 3-connected triangulations with a small number of separating triangles.

2. ONE SEPARATING TRIANGLE

The theorem which we will prove here is mainly a corollary of Theorem 1.1. We will show that for triangulations we can relax the restriction on 4-connectedness a bit and still have that the graph is hamiltonian-connected. More specific we will prove the following theorem.

Theorem 2.1. *Let G be a 3-connected triangulation. Let uv be an edge which is contained in all separating triangles of G . Then G is hamiltonian-connected.*

As an immediate and easy corollary of this theorem, we get that this property trivially also holds for triangulation with exactly one separating triangle.

Corollary 2.2. *Let G be a 3-connected triangulation with exactly one separating triangle. Then G is hamiltonian-connected.*

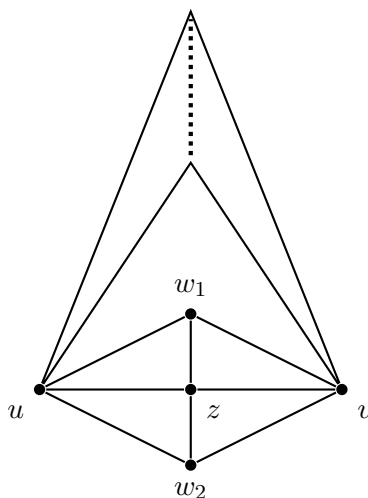


FIGURE 1. Subdividing the edge shared by all separating triangles.

Proof of Theorem 2.1. Let G be a 3-connected triangulation and let uv be an edge which is contained in all separating triangles of G . The edge uv is contained in two facial triangles of G which we denote by uvw_1 and uvw_2 . We obtain the graph G' by subdividing the edge uv with the vertex z and connecting z with the vertices w_1 and w_2 . The graph G' is shown in Figure 1. The modification adds a vertex in each separating triangle and no new separating triangles are created, so G' is 4-connected.

In order to show that G is hamiltonian-connected, we have to show for each pair of distinct vertices $x, y \in V(G)$ that there is a hamiltonian path from x to y . Owing to Theorem 1.2, we have that G' contains a hamiltonian path P' from x to y containing zv . The path P' contains at some point the sequence $*zv$ where $* \in \{w_1, w_2, u\}$ (or $* \in \{w_1, w_2\}$ if $xy = uv$). We obtain the path P from P' by replacing the sequence $*zv$ by the edge $*v$. Since all vertices and edges of G' except for z and its incident edges are also vertices and edges in G , we have that the path P is a hamiltonian path from x to y in G . \square

3. DECOMPOSITION TREES WITH MAXIMUM DEGREE $\Delta(T) > 3$

Jackson and Yu [4] defined a decomposition tree for a 3-connected triangulation as follows: A triangulation with a separating triangle S can be split into two triangulations: the subgraphs inside and outside of the separating triangle with a copy of S contained in both. By iteratively applying this procedure to a triangulation with k separating triangles, we obtain a collection of $k + 1$ triangulations without separating triangles. These triangulations form the vertices of the decomposition tree and we refer to them as the (4-connected) pieces of the original triangulation. Note that such

a ‘4-connected’ piece is either a 4-connected triangulation or K_4 . Jackson and Yu explain in [4] that this process indeed yields a tree and that this tree is unique. Two vertices are adjacent if the corresponding pieces share a separating triangle in the original triangulation.

Reversely it is easily seen that each tree can appear as the decomposition tree of some triangulation. The basic operation used to construct such a triangulation is that of subdivision. Given two triangulations G_1 and G_2 with no separating triangles and a (facial) triangle in both, we can subdivide G_1 with G_2 , or equivalently G_2 with G_1 , by identifying the two triangles (note that this is not unique) such that G_1 and G_2 are placed at different sides of the triangle. This process leads to a triangulation with one separating triangle and thus K_2 as decomposition tree and G_1 and G_2 as 4-connected pieces. Given a tree we can construct a triangulation with that tree as its decomposition tree by picking for each vertex of the tree an arbitrary triangulation without separating triangles with at least as many faces as the degree of the vertex of the tree, and then repeatedly subdivide the triangulations with each other if the corresponding vertices are adjacent in the tree.

In [5] the *scattering number* $s(G)$ of a graph G is defined as

$$s(G) = \max\{k(G - X) - |X| : X \subseteq V(G), k(G - X) \neq 1\},$$

where $k(H)$ denotes the number of components of the graph H . In [3] it is noted that $s(G) \leq -1$ is a necessary condition for G to be hamiltonian-connected.

The wheel graph W_n on n vertices is the plane graph obtained by adding a vertex in the center of a cycle of length n and connecting it to all vertices of the cycle.

Theorem 3.1. *For each tree T with maximum degree $\Delta(T) > 3$, there exists a 3-connected triangulation G with T as decomposition tree such that G is not hamiltonian-connected.*

Proof. Let v be a vertex of T with degree $d(v) > 3$. Removing v from T results in $d(v)$ components which we denote by $T_1, \dots, T_{d(v)}$. Using the technique described above, we can construct a triangulation G_i for each of the trees T_i such that G_i has T_i as a decomposition tree. At this point we have no requirements for each triangulation G_i except that it has T_i as decomposition tree and that the piece corresponding to the vertex of T_i that neighbored v in T has a facial triangle that is not subdivided. Let G be the graph obtained by subdividing each of the triangular faces of $W_{d(v)}$ with one of the G_i 's, and subdividing the outer face with a single vertex.

Since $d(v)$ is strictly greater than three, we have that G is a simple triangulation with decomposition tree T . Note that if $d(v)$ was equal to three, then G would still be a simple triangulation, but the last step would have added an extra separating triangle, so G would have had a decomposition tree different from T .

We have that $k(G - V(W_{d(v)})) = d(v) + 1$ and $|V(W_{d(v)})| = d(v) + 1$, so $s(G) \geq 0$. This implies that G is not hamiltonian-connected. \square

This technique cannot be used to exclude any subcubic tree as the decomposition tree of a hamiltonian-connected triangulation. In [2] it is shown that in a triangulation G with a subcubic decomposition tree, we have for each splitting set S : $k(G - S) < |S|$. So the scattering number of G is at most -1 .

4. COMPUTATIONAL RESULTS

In order to check whether a graph with n vertices is hamiltonian-connected, we need to check whether there exists a hamiltonian path between $n(n - 1)/2$ pairs of vertices. If two adjacent vertices in a graph on $n \geq 3$ vertices are connected by a hamiltonian path, then we actually have a hamiltonian cycle, and we can conclude that all vertices that are adjacent on the cycle are connected by a hamiltonian path. This can be used to eliminate several pairs while looking for hamiltonian paths. Below we will see some more results which can further speed up the programs that verify whether a triangulation is hamiltonian-connected.

We will need a theorem proven by Jackson and Yu in [4].

Theorem 4.1 ([4]). *Let G be a 3-connected triangulation with a decomposition tree T such that T has maximum degree at most three. Let H be a piece of G corresponding to a vertex of degree at most 2 in T , t be a facial cycle of both H and G , and $V(t) = \{u, v, w\}$. Then G has a hamiltonian cycle through uv and vw .*

The next corollary immediately follows from this theorem.

Corollary 4.2. *Let G be a 3-connected triangulation with a path as decomposition tree. Then G has a hamiltonian cycle through any edge.*

Proof. Let uv be an edge of G . Pick one of the two faces containing the edge uv , and label the third vertex w . Since the maximum degree of the decomposition tree is 2, all conditions for Theorem 4.1 are met, so G has a hamiltonian cycle through uv and vw , so certainly through uv . \square

If a triangulation has two separating triangles, then the decomposition tree is always a path of length 2. If it has three separating triangles, then the decomposition tree is either a path of length 3 or $K_{1,3}$.

The following lemma was useful for faster checking whether a graph is hamiltonian-connected, since it allows us to decide on the existence of several other hamiltonian paths based on a single hamiltonian path.

Lemma 4.3. *Let G be a graph with n vertices. Let P be a hamiltonian path in G . Let x_1, \dots, x_n be the sequence of vertices on this path (so in P we have that x_i is adjacent to x_{i+1} for $1 \leq i < n$). If x_i ($1 < i \leq n$) is adjacent to x_1 in G , then there is a hamiltonian path from x_{i-1} to x_n .*

Proof. The path from x_{i-1} to x_n is given by $x_{i-1} \dots x_1 x_i \dots x_n$. \square

The lemma above is obviously more powerful if the degree of the end points of the paths are large. Therefore we sorted the vertices, so that we first checked for paths between the vertices with the largest degrees. After applying this lemma, we get several hamiltonian paths between new pairs of vertices. For each of these paths, we can again apply this lemma. This also vastly increased the speed of the program. Applying the lemma a third time to the new paths did not really deliver a significant increase in speed.

The following two lemmata do not give any information on the existence of a hamiltonian path between two specific vertices, but reduces the nonexistence of such a path to the nonexistence of another path in a different triangulation.

Lemma 4.4. *Let G be a triangulation on n vertices containing s separating triangles. Let (u, v, w_1) and (u, v, w_2) be two facial triangles of G such that w_1 and w_2 are not adjacent and $N(w_1) \cap N(w_2) = \{u, v\}$. If each pair of adjacent vertices in all triangulations on n vertices containing at most s separating triangles are connected by a hamiltonian path, then G has a hamiltonian path from w_1 to w_2 .*

Proof. Consider the graph G' obtained from G by removing the edge uv and adding the edge w_1w_2 . Since w_1 and w_2 are not adjacent in G , we have that G' is a simple triangulation. The intersection of the neighbourhoods of w_1 and w_2 is $\{u, v\}$, so G' has at most s separating cycles. Together we find that the graph G' is a triangulation on n vertices having at most s separating cycles, so it contains a hamiltonian path P from w_1 to w_2 . Since all edges of G' except for w_1w_2 are also contained in G , we have that P is also a hamiltonian path from w_1 to w_2 in G . \square

An edge in a 3-connected simple triangulation is called *reducible* if it is not contained in a separating triangle or a chordless separating quadrangle. The following lemma allows us to skip some pairs of adjacent vertices.

Lemma 4.5. *Let G be a triangulation on n vertices with a decomposition tree D . Let u be a vertex of degree 4, and let uv be a reducible edge of G . If each pair of adjacent vertices in all triangulations on $n - 1$ vertices with decomposition tree D are connected by a hamiltonian path, then G has a hamiltonian path from u to v .*

Proof. Since the edge uv is not contained in a separating triangle or a chordless separating quadrangle, the graph obtained by contracting uv in G will still have the same decomposition tree, even if uv is contained in separating quadrangle with a chord. Let v, w_1, w_2, w_3 be the cyclic order of the vertices around u . Let G' be the triangulation obtained from G by removing u and its incident edges and adding the edge vw_2 , i.e., G' is the triangulation obtained by contracting the edge uv . In G' , there is a hamiltonian path P

from v to w_2 . All edges in G' except the edge vw_2 are also contained in G , so $P \cup \{uw_2\}$ is a hamiltonian path from u to v in G . \square

The lemma above gives no new information for triangulations where the decomposition tree is a path, since in that case we already know that adjacent vertices are connected by a hamiltonian path. The following two lemmata prove similar results for triangulations with a path as decomposition tree, but this time for certain vertices at distance two.

Lemma 4.6. *Let G be a triangulation with a path as decomposition tree. Let u be a vertex of degree 4, and let v_1, v_2, v_3, v_4 be the cyclic order of the vertices around u . Let uv_1 be a reducible edge of G . Then G has a hamiltonian path from v_1 to v_3 .*

Proof. Similar to the proof in the previous lemma, we can conclude that the triangulation G' obtained by contracting the edge uv_1 in G has a path as decomposition tree. Owing to Theorem 4.1, G' has a hamiltonian cycle C through v_1v_2 and v_1v_3 . A hamiltonian path from v_1 to v_3 in G is $(C \setminus \{v_1v_2, v_1v_3\}) \cup \{uv_1, uv_2\}$. \square

Lemma 4.7. *Let G be a triangulation with a path as decomposition tree. Let u be a vertex of degree 5, and let v_1, v_2, v_3, v_4, v_5 be the cyclic order of the vertices around u . Let uv_1 be a reducible edge of G . Then G has a hamiltonian path from v_1 to v_3 , and from v_1 to v_4 .*

Proof. We will only give the proof for the hamiltonian path from v_1 to v_3 . The other proof is completely analogous.

Similar to the proof in the previous lemma, we can conclude that the triangulation G' obtained by contracting the edge uv_1 in G has a path as decomposition tree. Owing to Theorem 4.1, G' has a hamiltonian cycle C through v_1v_3 and v_3v_4 . A hamiltonian path from v_1 to v_3 in G is $(C \setminus \{v_1v_3, v_3v_4\}) \cup \{uv_3, uv_4\}$. \square

The program described above was compared with an independent implementation which checks for each pair of vertices whether a hamiltonian cycle exists in the graph obtained by adding a single vertex and connecting it to the two vertices in the pair. It was tested for 3-connected triangulations how many are hamiltonian-connected. Owing to Corollary 4.2, only pairs of nonadjacent vertices need to be checked in case of two separating triangles, but no further optimisations were used in this implementation for testing purposes.

The program with the optimisations was used to test 3-connected triangulations with two separating triangles up to 22 vertices. For 22 vertices it was run on a cluster of Intel Xeon E5-2680 CPU's at 2.5 GHz. There were 21 282 658 291 triangulations with two separating triangles and the computations took about 22.7 CPU years. The result of this computation is the following lemma:

Lemma 4.8. *On up to 22 vertices all 3-connected triangulations with two separating triangles are hamiltonian-connected.*

The program with the optimisations was also used to test 3-connected triangulations with three separating triangles up to 21 vertices. For 21 vertices it was run on a cluster of Intel Xeon E5-2680 CPU's at 2.5 GHz. There were 8 751 268 952 triangulations with three separating triangles and the computations took about 6.3 CPU years. The result of this computation is the following lemma:

Lemma 4.9. *On up to 21 vertices all 3-connected triangulations with three separating triangles are hamiltonian-connected.*

We also tested 3-connected triangulations with more than 3 separating triangles, but with specific decomposition trees. We verified hamiltonian-connectedness for triangulations with a path as decomposition tree up to 21 vertices. For 21 vertices it was run on a cluster of Intel Xeon E5-2680 CPU's at 2.5 GHz. There were 10 141 293 048 triangulations with a path as a decomposition tree and the computations took about 5.4 CPU years. The result of this computation is the following lemma:

Lemma 4.10. *On up to 21 vertices all 3-connected triangulations with a path as decomposition tree are hamiltonian-connected.*

We also verified hamiltonian-connectedness for triangulations with a decomposition tree with maximum degree 3 up to 20 vertices. For 20 vertices it was run on a cluster of Intel Xeon E5-2680 CPU's at 2.5 GHz. There were 23 748 083 814 triangulations with a decomposition tree with maximum degree 3 and the computations took about 8.0 CPU years. The result of this computation is the following lemma:

Lemma 4.11. *On up to 20 vertices all 3-connected triangulations with a decomposition tree with maximum degree 3 are hamiltonian-connected.*

5. CONCLUSION

It is known that 4-connected triangulations are hamiltonian-connected. We showed that also 3-connected triangulations with only one separating triangle are hamiltonian-connected. If we express this in terms of the decomposition tree, then we have that 3-connected triangulations with a decomposition tree with maximum degree at most 1 are hamiltonian-connected. In order to show bounds on the strongest form of this theorem, we proved that any decomposition tree with maximum degree at least 4 can appear as the decomposition tree of a 3-connected triangulation that is not hamiltonian-connected. This also means that if $s \geq 4$ then there is always a 3-connected triangulation with s separating triangles that is not hamiltonian-connected. Finally, we presented some computational results which show that all 'small' 3-connected triangulations with a decomposition tree with maximum degree at most 3 are hamiltonian-connected.

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