

**BOUNDS ON THE ACHROMATIC NUMBER OF PARTIAL  
TRIPLE SYSTEMS**

PETER DUKES, GARY MACGILLIVRAY, AND KRISTIN PARTON

ABSTRACT. A complete  $k$ -colouring of a hypergraph is an assignment of  $k$  colours to the points such that (1) there is no monochromatic hyperedge, and (2) identifying any two colours produces a monochromatic hyperedge. The achromatic number of a hypergraph is the maximum  $k$  such that it admits a complete  $k$ -colouring. We determine the maximum possible achromatic number among all maximal partial triple systems, give bounds on the maximum and minimum achromatic numbers of Steiner triple systems, and present a possible connection between optimal complete colourings and projective dimension.

## 1. INTRODUCTION

A  $t$ -uniform hypergraph is a pair  $(V, \mathcal{A})$ , where  $V$  is a nonempty set of points and  $\mathcal{A}$  is a collection of  $t$ -subsets of  $V$ . Here, elements of  $\mathcal{A}$  are called blocks. A partial triple system of order  $v$ , abbreviated  $\text{PTS}(v)$ , is a 3-uniform hypergraph  $(V, \mathcal{B})$  with  $|V| = v$  and for which every pair of distinct elements of  $V$  is contained in at most one block. If every pair of distinct points in  $V$  occurs in exactly one block, then  $(V, \mathcal{B})$  is a Steiner triple system, or  $\text{STS}(v)$ . It is well-known that an  $\text{STS}(v)$  exists if and only if  $v$  is a positive integer with  $v \equiv 1, 3 \pmod{6}$ . In the case  $v = 1$ , we take  $\mathcal{B} = \emptyset$  to satisfy the conditions.

The leave of a  $\text{PTS}(v)$   $(V, \mathcal{B})$  is the graph  $(V, E)$  with  $xy$  an edge if and only if  $x$  and  $y$  are together in no block of  $\mathcal{B}$ . A  $\text{PTS}(v)$  is maximal if its leave is triangle-free. A consequence of Mantel's famous theorem on triangle-free graphs is that a maximal  $\text{PTS}(v)$  has at least  $v(v-2)/12$  blocks. A  $\text{PTS}(u)$   $(U, \mathcal{A})$  is a subsystem of (or embeds in) a  $\text{PTS}(v)$   $(V, \mathcal{B})$  if  $U \subseteq V$  and  $\mathcal{A} \subseteq \mathcal{B}$ . For later use, we state a recent result on embedding  $\text{PTS}$  into  $\text{STS}$ .

**Lemma 1.1** ([1]). *Suppose  $v > 2u$  and  $v \equiv 1, 3 \pmod{6}$ . Then every  $\text{PTS}(u)$  is a subsystem of some  $\text{STS}(v)$ .*

---

Received by the editors November 1, 2005, and in revised form August 7, 2006.  
2000 *Mathematics Subject Classification.* 05B07, 05C15.

*Key words and phrases.* achromatic number, complete colouring, hypergraph, partial triple system, Steiner triple system.

Research of the authors is supported by NSERC.

A *proper  $k$ -colouring* of a  $t$ -uniform hypergraph  $(V, \mathcal{A})$  is a mapping  $c : V \rightarrow K$ , where  $|K| = k$ , such that  $|c(A)| > 1$  for each  $A \in \mathcal{A}$ . The elements of  $K$  are called *colours*, and we may assume that  $K = \{1, \dots, k\}$ . The  $i$ th *colour class* is  $c^{-1}(i)$ , the set of points assigned colour  $i$ .

A proper  $k$ -colouring  $c$  is *complete* if for every pair  $i, j$  of distinct colours there is a block  $A \in \mathcal{A}$  with  $c(A) = \{i, j\}$ . We say that the pair  $\{i, j\}$  is *covered by  $A$* . The *achromatic number* of  $(V, \mathcal{A})$ , denoted  $\psi(\mathcal{A})$ , is the maximum  $k$  such that  $(V, \mathcal{A})$  admits a complete  $k$ -colouring.

For each positive integer  $n$ , define  $\psi_{\min}(v)$  and  $\psi_{\max}(v)$  as the minimum and maximum achromatic numbers, respectively, of a maximal PTS( $v$ ). For  $v \equiv 1, 3 \pmod{6}$ , define  $\psi_{\min}^*(v)$  and  $\psi_{\max}^*(v)$  similarly for STS( $v$ ). When meaningful, it is clear that

$$\psi_{\min}(v) \leq \psi_{\min}^*(v) \leq \psi_{\max}^*(v) \leq \psi_{\max}(v).$$

In the next section, we give an explicit upper bound on  $\psi_{\max}(v)$ . Then, in section 3, we show that this upper bound is met with equality by constructing a PTS( $v$ ) with this achromatic number. Bounds for  $\psi_{\max}^*(v)$  follow as a consequence. In section 4, we modify an argument in [3] to establish lower bounds on  $\psi_{\min}(v)$  and  $\psi_{\min}^*(v)$ . We present upper bounds and open problems concerning these quantities in section 5. Finally, various optimal complete colourings of small STS are compiled in section 6.

## 2. UPPER BOUNDS ON $\psi_{\max}$

In [3], it was shown that  $\psi_{\max}(v)$  is  $O(v^{2/3})$ . With straightforward counting, we are able to obtain an exact upper bound.

**Lemma 2.1.** *If there exists a complete  $k$ -colouring of a PTS( $v$ ) with colour class sizes  $y_1 \leq y_2 \leq \dots \leq y_k$ , then  $\sum_{i=1}^j \binom{y_i}{2} \geq \binom{j}{2}$  for all  $j = 1, \dots, k$ .*

*Proof.* Given a complete  $k$ -colouring  $c : V \rightarrow K$  of  $(V, \mathcal{B})$ , define a digraph  $D$  with vertex set  $K$  and with  $(i, j)$  an arc if and only if there exists a block  $B \in \mathcal{B}$  with  $|c^{-1}(i) \cap B| = 2$  and  $|c^{-1}(j) \cap B| = 1$  (that is, if  $B$  has two points coloured  $i$  and one point coloured  $j$ ). Let  $s_1 \leq s_2 \leq \dots \leq s_k$  be the sequence of outdegrees of  $D$ . We have  $s_i \leq \binom{y_i}{2}$ . By completeness,  $D$  (and each of its induced subdigraphs) is semi-complete. So  $\sum_{i=1}^j \binom{y_i}{2} \geq \sum_{i=1}^j s_i \geq \binom{j}{2}$  for  $1 \leq j \leq k$ .  $\square$

**Corollary 2.2.** *In a complete  $k$ -colouring of a PTS( $v$ ), there are at most  $t^2 - t + 1$  colour classes of size  $\leq t$ .*

*Proof.* Suppose there are  $k$  colour classes of size  $\leq t$ . By Lemma 2.1,  $k \binom{t}{2} \geq \binom{k}{2}$ , from which the result follows.  $\square$

**Remark 2.3.** *These results can be generalized for triple systems of higher index  $\lambda$  in which every pair of distinct points belongs to at most  $\lambda$  blocks.*

For  $n \geq 1$ , define  $a_n = \lceil \sqrt{2n-1} \rceil$ . The sequence  $\{a_n\}$  begins

$$1, 2, 3, 3, 3, 4, 4, 4, 5, 5, 5, 5, 5, 6, 6, 6, 6, 6, \dots, \overbrace{2i-1, \dots, 2i-1}^{2i-1}, \overbrace{2i, \dots, 2i}^{2i-1}, \dots$$

The following facts are easily verified by induction. We omit the proofs.

- (1)  $\sum_{i=1}^j \binom{a_i}{2} \geq \binom{j}{2}$  for all  $j$ , with equality if and only if  $j = 2n^2$  or  $j = 2n^2 + 2n + 1$  for some  $n$ ;
- (2) if, for all  $j = 1, \dots, k$ , we have  $b_j \leq a_j$  and  $\sum_{i=1}^j \binom{b_i}{2} \geq \binom{j}{2}$ , then  $b_j = a_j$  for all  $j = 1, \dots, k$ .

Taken with Lemma 2.1, (1) and (2) lead to an upper bound on  $\psi_{\max}$ .

**Theorem 2.4.** *Let  $a_n = \lceil \sqrt{2n-1} \rceil$ . Then*

$$\psi_{\max}(v) \leq \max\{k : a_1 + a_2 + \dots + a_k \leq v\}.$$

*Proof.* Let  $k$  denote the right side above, and assume there is a complete  $(k+1)$ -colouring of some  $\text{PTS}(v)$ , say with colour class sizes  $y_1 \leq y_2 \leq \dots \leq y_{k+1}$ . By Lemma 2.1, we have  $\sum_{i=1}^j \binom{y_i}{2} \geq \binom{j}{2}$  for  $j = 1, \dots, k+1$ . We repeatedly transform the sequence of  $y_i$  as follows. Since  $\sum_{i=1}^{k+1} y_i < \sum_{i=1}^{k+1} a_i$ , we may choose the smallest integer  $m \in \{1, \dots, k+1\}$  such that  $y_m < a_m$ . By (2) above, we cannot have  $y_i = a_i$  for all  $i < m$ , so take the largest integer  $n \in \{1, \dots, m-1\}$  with  $y_n > a_n$ . Now define integers  $w_1, \dots, w_{k+1}$  with  $w_n = y_n - 1$ ,  $w_m = y_m + 1$ , and  $w_i = y_i$  for all  $i \neq n, m$ . Now property (1) above implies

$$\sum_{i=1}^j \binom{w_i}{2} \geq \sum_{i=1}^j \binom{a_i}{2} \geq \binom{j}{2}$$

for  $j < m$  and since  $y_m \geq y_n > w_n$ , we have

$$\sum_{i=1}^j \binom{w_i}{2} = \sum_{i=1}^j \binom{y_i}{2} + (y_m - w_n) > \binom{j}{2}$$

for  $j \geq m$ . Now relabel each  $w_i$  as  $y'_i$  and choose indices  $m'$  and  $n'$  for  $y'_i$  as  $m$  and  $n$  were chosen for  $y_i$ . It is noteworthy that  $y'_{m'} \geq y_{m'} \geq y_{n'} \geq y'_{n'}$ , so the estimate on  $\sum_{i=1}^j \binom{w_i}{2}$  for  $j \geq m$  remains valid in subsequent steps. This process must terminate, contradicting (2). Therefore, the supposed colouring does not exist.  $\square$

It is particularly interesting to consider the case of equality in (1). For  $k = 2n^2$ , we calculate

$$\sum_{i=1}^k a_i = \sum_{i=1}^n (2i-1)(4i-1) = \frac{8}{3}n^3 + n^2 - \frac{2}{3}n,$$

and similarly, for  $k = n^2 + (n+1)^2$ ,

$$\sum_{i=1}^k a_i = \frac{8}{3}n^3 + 5n^2 + \frac{10}{3}n + 1.$$

For future reference, we label these cubic polynomials in  $n$  as  $p_1(n)$  and  $p_2(n)$ , respectively. The following consequence of (1), (2), and Theorem 2.4 is now immediate.

**Corollary 2.5.** *Let  $k = 2n^2$ , (or  $k = n^2 + (n + 1)^2$ ), where  $n$  is a positive integer. If there is a complete  $k$ -colouring of some  $\text{PTS}(v)$ , then  $v \geq p_1(n)$  (respectively  $v \geq p_2(n)$ ), with equality if and only if the colour class sizes are exactly  $a_1, a_2, \dots, a_k$ .*

By Corollary 2.5 and condition (1) on the sequence  $\{a_n\}$ , when  $v = p_1(n)$  or  $v = p_2(n)$ , any  $\psi_{\max}(v)$ -colouring of a  $\text{PTS}(v)$  uses every pair of points within a colour class to cover some pair of colours, and that each pair of colours is covered exactly once. We consider an application of this structure in section 5.

### 3. LOWER BOUNDS ON $\psi_{\max}$ AND $\psi_{\max}^*$

We begin with a construction of a family of PTS having largest possible achromatic number.

**Theorem 3.1.** *Let  $a_n = \lceil \sqrt{2n-1} \rceil$ . Then*

$$\psi_{\max}(v) \geq \max\{k : a_1 + a_2 + \dots + a_k \leq v\}.$$

*Proof.* For  $h > 1$ , define  $[h] = \max\{l : a_l < a_h\}$ . We construct a PTS on points  $X_1 \cup \dots \cup X_k$ , where the  $X_i$  are pairwise disjoint colour classes with  $|X_i| = a_i$ , and every pair of colours is covered by some block. It is sufficient to perform this construction for an infinite sequence of  $k$ . Hence, we will assume  $[k+1] = k$ , by replacing  $k$  by the least integer  $K$  with  $[K] > [k]$ , and deleting if necessary the points in  $X_{k+1} \cup \dots \cup X_K$ . For  $j > 1$ , order the pairs in  $X_j$  arbitrarily. Form blocks by joining the  $i$ th pair in  $X_j$  to some point in  $X_i$ , for each  $i = 1, \dots, [j]$ . Note that

$$\binom{a_j}{2} - [j] = \lceil a_j/2 \rceil - 1,$$

so every pair of colours  $i < j$  with  $a_i < a_j$  is now covered by some block. For a given  $j$  with  $1 < j \leq k$ , there are  $b_j = 2\lceil a_j/2 \rceil - 1$  values of  $i$  with  $[i] = [j]$ , or  $a_i = a_j$ . It remains to define blocks covering the pairs  $i, j$ , with  $[i] < i < j \leq [i] + b_j$ . Consider the complete graph  $K_{b_j}$  on vertices  $1, \dots, b_j$ . Orient its edges such that every vertex has indegree and outdegree equal to  $\lceil a_j/2 \rceil - 1$ . (For instance, this can be done using a decomposition into Hamilton cycles.) For each edge directed from  $r$  to  $s$ , form a block by joining a unique choice of one of the remaining  $\binom{a_j}{2} - [j]$  pairs of  $X_{[i]+r}$  to some point in  $X_{[i]+s}$ .  $\square$

**Corollary 3.2.**

$$\psi_{\max}(v) = \max\{k : a_1 + a_2 + \dots + a_k \leq v\}.$$

*Proof.* This follows directly from Theorems 2.4 and 3.1.  $\square$

The asymptotic behavior of  $\psi_{\max}$  is easily calculated from the definition of  $a_n$ .

**Corollary 3.3.**

$$\lim_{v \rightarrow \infty} \frac{\psi_{\max}(v)}{v^{2/3}} = \frac{1}{2} \cdot 3^{2/3}.$$

We now turn to the question of lower bounds on  $\psi_{\max}^*(v)$ , the maximum achromatic number of a Steiner triple system of order  $v$ .

Suppose a PTS  $(U, \mathcal{A})$  is a subsystem of another PTS  $(V, \mathcal{B})$ . Let us call a colouring of  $(U, \mathcal{A})$  *safe with respect to  $\mathcal{B}$*  if it induces no monochromatic block in  $\mathcal{B}$ .

**Lemma 3.4.** *Suppose  $(U, \mathcal{A})$  is a PTS( $u$ ) with a complete  $k$ -colouring that is safe with respect to some  $\mathcal{B} \supseteq \mathcal{A}$ . Then a PTS( $v$ )  $(V, \mathcal{B})$  has a complete  $l$ -colouring for some  $l \geq k$ .*

*Proof.* Consider a PTS( $v$ )  $(V, \mathcal{B})$  with  $U \subseteq V$ . Initially colour the points of  $U$  with a complete  $k$ -colouring, safe with respect to  $\mathcal{B}$ , using colours  $1, \dots, k$ , and the points of  $V \setminus U$  each with a distinct new colour, using colours  $k + 1, \dots, k + v - u$ . Since the colouring of  $U$  is safe, it follows that this colouring of  $V$  is proper. Now repeatedly perform the following operation: merge any two colour classes  $i < j$  which are not covered by any block, and rename this colour class  $i$ . It is clear that classes  $1, \dots, k$  remain nonempty, and when no merging is possible, a complete colouring results.  $\square$

It is evident that the construction in Theorem 3.1 can be done in such a way that the resulting colouring is safe with respect to any  $\mathcal{B} \supseteq \mathcal{A}$  (for instance, if the  $u$  “unused” pairs of points in a colour class are chosen to form a star  $K_{1,u}$ ). Together with Lemmas 1.1 and 3.4, we obtain the following result.

**Theorem 3.5.**

$$\psi_{\max}^*(v) \geq \max\{k : a_1 + a_2 + \dots + a_k < v/2\}.$$

**Corollary 3.6.** *If  $\lim_{v \rightarrow \infty} \psi_{\max}^*(v)/v^{2/3}$  exists and equals  $L$ , then*

$$\frac{1}{2} \cdot (3/2)^{2/3} \leq L \leq \frac{1}{2} \cdot 3^{2/3}.$$

In practice, it seems easy to embed *some* PTS( $v$ ) from Theorem 3.1 into *some* STS( $v$ ), provided  $v \equiv 1, 3 \pmod{6}$ . Using a standard hill-climbing algorithm for completing STS, we were able to do so for all “small” values of  $v$ . Some specific constructions are given in the appendix.

**Theorem 3.7.** *If  $v \equiv 1, 3 \pmod{6}$  and  $1 \leq v \leq 49$ , then  $\psi_{\max}^*(v) = \psi_{\max}(v)$ .*

An interesting question is whether such an embedding is always possible. If so, the following would be proved.

**Conjecture 3.8.** *For all  $v \equiv 1, 3 \pmod{6}$ ,  $\psi_{\max}^*(v) = \psi_{\max}(v)$ .*

4. LOWER BOUNDS ON  $\psi_{\min}$  AND  $\psi_{\min}^*$ 

There is a unique STS( $v$ ) up to isomorphism for  $v \leq 9$ . So  $\psi_{\min}^*(v) = \psi_{\max}^*(v)$  for  $v \leq 9$ . For  $v = 13, 15$ , there are, respectively, 2 and 80 different STS up to isomorphism. After a very fast computer search, we report that every STS(13) and STS(15) admits a complete 5-colouring.

**Theorem 4.1.**  $\psi_{\min}^*(v) = 5$  for  $v = 13, 15$ .

The following is adapted from [3].

**Theorem 4.2.** Any PTS with minimum degree  $t \geq 3$  admits a complete  $k$ -colouring, where

$$k = \left\lfloor \frac{11 + \sqrt{12t - 11}}{6} \right\rfloor.$$

*Proof.* Let  $(V, \mathcal{B})$  be such a PTS( $v$ ). Since  $t \geq 3$ , we have  $2 \leq k < t$ . Take  $x_1 \in V$  and blocks  $B_1, B_2, \dots, B_{k-1}$ , each containing  $x_1$ . Define  $X_1 = \{x_1\}$ ,  $X_{1,i} = B_i \setminus \{x_1\}$  for each  $i = 1, 2, \dots, k-1$ . Suppose for some  $r \geq 1$  we have constructed pairwise disjoint sets

$$X_1, X_2, \dots, X_r, X_{r,r}, X_{r,r+1}, \dots, X_{r,k-1} \subset V,$$

where  $|X_1| = 1, |X_i| = 2i - 2$  for  $i = 2, 3, \dots, r$ , and  $|X_{r,j}| = 2r$  for  $j = r, r+1, \dots, k-1$ . Further suppose that

- (1) for any  $i, j \leq r$  with  $i \neq j$ ,  $X_i \cup X_j$  contains a block of  $\mathcal{B}$ ;
- (2) for any  $i \leq r$  and  $j \geq r$ ,  $X_i \cup X_{r,j}$  contains a block of  $\mathcal{B}$ ; and
- (3) there does not exist  $i$  such that  $X_i$  or  $X_{r,i}$  contains a block of  $\mathcal{B}$ .

Evidently, these conditions hold for  $r = 1$ . If  $r = k-1$ , define  $X_k = X_{k-1, k-1}$  and the construction is done. The sets  $X_1, \dots, X_k$  are colour classes of a complete  $k$ -colouring of some subsystem  $(U, \mathcal{A})$  which is safe with respect to  $\mathcal{B}$ . Lemma 3.4 extends this colouring to  $(V, \mathcal{B})$ . It remains to show that for  $r < k-1$  we can continue the construction.

Take a point  $x_{r+1} \in X_{r,r}$  and define

$$W_r = \left( \bigcup_{i=1}^r X_i \right) \cup \left( \bigcup_{j=r}^{k-1} X_{r,j} \right),$$

the set of all points used in the construction so far. There are at most  $|W_r| - 1 = r(r-1) + 2r(k-r)$  blocks of the form  $\{x_{r+1}, y, w\}$  with  $w \in W_r$ .

Now for  $i = r+1, \dots, k-1$ , let  $Z_{r,i}$  be the set of points  $z \notin W_r$  such that for some  $B \in \mathcal{B}$ ,  $z \in B$  and  $B \setminus \{z\} \subseteq X_{r,i}$ . For such values of  $i$ , we are forbidden by (3) above to extend  $X_{r,i}$  to  $X_{r+1,i}$  by adding a point in  $Z_{r,i}$ . Note that  $|Z_{r,i}| \leq \binom{2r}{2} - r = 2r(r-1)$  for each  $i$ , since there are this many ‘‘unused’’ pairs in  $X_{r,i}$ .

It follows that there are at least  $t - 3r(r-1) - 2r(k-r)$  choices for a block  $B_i$  containing  $x_{r+1}$  with  $B_i \setminus x_{r+1}$  disjoint from  $W_r$  and extending  $X_{r,i}$

without violating (3). Since we must pick such blocks for  $i = r + 1, \dots, k - 1$ , the construction can continue if and only if

$$t - 3r(r - 1) - 2r(k - r) \geq k - r - 1,$$

or equivalently

$$r^2 + 2kr - 4r + k - 1 \leq t.$$

The quadratic in  $r$  on the left is minimized for  $r = 2 - k$ . For our purposes,  $1 \leq r < k - 1$ , so the left side is maximized for  $r = k - 2$ . The condition becomes  $3k^2 - 11k + 11 \leq t$ , or

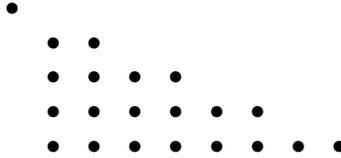
$$k \leq \frac{11 + \sqrt{12t - 11}}{6}.$$

Having chosen distinct blocks  $B_i \in \mathcal{B}$  which contain  $x_{r+1}$  and no other point of  $W_r \cup Z_{r,i}$ , we put

$$X_{r+1} = X_{r,r} \text{ and } X_{r+1,i} = X_{r,i} \cup (B_i \setminus \{x_{r+1}\})$$

for  $i = r + 1, \dots, k - 1$ . Observe that properties (1), (2), and (3) above hold for  $X_1, \dots, X_{r+1}, X_{r+1,r+1}, \dots, X_{r+1,k-1}$ .  $\square$

We remark that the argument proving Theorem 4.2 can be slightly improved with a careful choice of points and colour classes. We omit the details here. The construction of  $X_1, \dots, X_k$  is illustrated for  $k = 5$  below. The  $i$ th row represents  $X_i$ , with the rightmost point representing a choice of  $x_i$ . Blocks joining this point to pairs across the next two columns are present in  $\mathcal{B}$ .



**Corollary 4.3.** For  $v \geq 7$ ,

$$\psi_{\min}^*(v) \geq \frac{11 + \sqrt{6v - 17}}{6}.$$

*Proof.* Let  $t = \frac{v-1}{2}$  in Theorem 4.2.  $\square$

**Corollary 4.4.** Suppose  $v > 24$ . Then

$$\psi_{\min}(v) \geq \frac{11 + \sqrt{v - 11}}{6}.$$

*Proof.* It is enough to show that any maximal  $\text{PTS}(v)$   $(V, \mathcal{B})$  with  $v > 24$ , has a sub-PTS with the minimum degree of any point at least  $v/12$ . Choose  $x \in V$  with degree  $< v/12$ , and delete it together with all blocks through it. Repeat this process until all points have degree  $\geq v/12$ . We claim that this process terminates with at least  $(v - 5)/2$  points left. Suppose  $\lfloor (v + 5)/2 \rfloor$  points, each of degree  $< v/12$ , have been removed. Then at least  $v(v - 2)/12 - v(v + 5)/24 = v(v - 9)/24$  blocks remain. But there can be

no more than  $\frac{1}{3}\binom{v-4}{2} = (v-4)(v-6)/24$  blocks left. This contradicts  $v > 24$ .  $\square$

### 5. UPPER BOUNDS ON $\psi_{\min}$

It is natural to ask whether any maximal PTS( $v$ ) has achromatic number less than  $\psi_{\max}(v)$ . The following gives an infinite family of such PTS.

**Theorem 5.1.** *Suppose  $v = p_1(n) = \frac{8}{3}n^3 + n^2 - \frac{2}{3}n$  or  $v = p_2(n) = \frac{8}{3}n^3 + 5n^2 + \frac{10}{3}n + 1$ , where  $n \geq 2$  is an integer. Then  $\psi_{\min}(v) < \psi_{\max}(v)$ .*

*Proof.* Assume first that  $v$  is even. Since  $v \geq 24$ , we can write  $v = u + u'$ , where  $u, u' \equiv 1$  or  $3 \pmod{6}$  and  $u, u' \geq 9$ . Define a PTS( $v$ )  $(V, \mathcal{B})$ , where  $\mathcal{B}$  is the union of blocks of an STS( $u$ ) and an STS( $u'$ ) on points  $U, U'$ , where  $U \cap U' = \emptyset$ . It is clear that  $(V, \mathcal{B})$  is a maximal PTS( $v$ ). If  $(V, \mathcal{B})$  were to admit a complete  $\psi_{\max}(v)$ -colouring, then by the discussion following Corollary 2.5, every colour class is either completely in  $U$  or completely in  $U'$ . Since  $v \geq 24$ , there is certainly a pair of colours uncovered by  $\mathcal{B}$ . This is a contradiction, and therefore  $\psi_{\min}(v) < \psi_{\max}(v)$ . The case when  $v(\geq 49)$  is odd is similar, except we write  $v = u + u' - 1$  and  $V = U \cup U'$ , where  $|U \cap U'| = 1$ .  $\square$

The question of whether  $\psi_{\min}^*(v) < \psi_{\max}^*(v)$  for any  $v$  seems more difficult. One approach would be to attempt a construction of STS( $v$ ) avoiding certain configurations of blocks required in an optimal colouring. Although we do not have an example of an STS( $v$ ) with a provably “bad” achromatic number, there is an infinite family of STS( $v$ ) for which we can deduce information about optimal colourings.

Given a PTS( $v$ ), say  $(V, \mathcal{B})$ , we say  $I \subset V$  is an *independent set* if there is no  $B \in \mathcal{B}$  with  $B \subseteq I$ . A complete  $k$ -colouring of a PTS( $v$ ) is equivalent to a partition  $I_1, \dots, I_k$  of  $V$  into independent sets such that  $I_i \cup I_j$  is not independent for  $i \neq j$ . From this observation and the pigeonhole principle, we obtain a structural result on colour class sizes.

**Theorem 5.2.** *Let  $\mathcal{I} = \{I_1, \dots, I_N\}$  be a family of independent sets of a PTS( $v$ ). Suppose for every positive integer  $i$  that each independent  $i$ -subset of  $V$  is contained in at least  $m_i$  elements of  $\mathcal{I}$ . In any complete  $k$ -colouring with  $n_i$  colour classes of size  $i$ , we have*

$$\sum_{i \geq 1} m_i n_i \leq N.$$

Let  $(V, \mathcal{B})$  be a fixed STS( $v$ ). Any STS( $(v-1)/2$ ),  $(U, \mathcal{A})$ , which is a subsystem is called a *projective hyperplane*. Note that for such  $U$ ,  $V \setminus U$  is necessarily a maximal (in fact, maximum) independent set.

For  $X \subset V$ , let  $\mathbf{e}_X : V \rightarrow \{0, 1\}$  denote the *characteristic vector* of  $X$ , where  $\mathbf{e}_X(x) = 1$  if and only if  $x \in X$ . The following is an observation of Teirlinck.

**Theorem 5.3** ([4]). *Let  $\mathcal{I}$  be the set of all  $I \subseteq V$ , in which  $V \setminus I$  induces a projective hyperplane of  $(V, \mathcal{B})$ . Then  $\mathcal{W} = \{\mathbf{e}_I : I \in \mathcal{I}\} \cup \{\mathbf{0}\}$  is a vector space over  $\mathbb{F}_2$ .*

Throughout, let  $d$  denote the dimension of  $\mathcal{W}$ . A well-known example of an STS whose independent sets induce a vector space of dimension  $d$  is now given. Let  $V = \mathbb{F}_2^{d+1} \setminus \{\mathbf{0}\}$ , the set of all nonzero binary  $(d+1)$ -tuples. For  $\mathbf{x}, \mathbf{y}, \mathbf{z} \in V$ , define  $\{\mathbf{x}, \mathbf{y}, \mathbf{z}\} \in \mathcal{B}$  if and only if  $\mathbf{x} + \mathbf{y} = \mathbf{z}$ . It is easy to see that  $(V, \mathcal{B})$  is an STS( $v$ ), where  $v = 2^{d+1} - 1$ , called the *projective STS of dimension  $d$* .

The following facts are easily verified with linear algebra.

**Lemma 5.4.** *In the projective STS of dimension  $d$ , any independent  $t$ -subset  $T$  of points of  $V$ ,  $0 \leq t \leq d+1$ , is contained in exactly  $2^{d+1-t} - 1$  projective hyperplanes, and is disjoint from exactly  $2^{d+1-t}$  projective hyperplanes for  $t \geq 1$ .*

By Theorem 5.2, we can now make a statement about colourings of projective STS.

**Corollary 5.5.** *Suppose there exists a complete colouring of the projective STS of dimension  $d$  with  $n_i$  colour classes of size  $i$ ,  $i = 1, \dots, d+1$ . Then*

$$\sum_{i=1}^{d+1} 2^{d+1-i} n_i \leq 2^{d+1} - 1.$$

It is perhaps unfortunate that the Diophantine equations  $p_1(n) = 2^d - 1$  and  $p_2(n) = 2^d - 1$  each have no solutions, as Corollaries 2.5 and 5.5 would lead to the conclusion that certain projective STS( $v$ ) have “bad” achromatic numbers. It appears that projective dimension is worth further attention in future work on  $\psi_{\min}^*(v)$ .

## 6. APPENDIX: COMPLETE COLOURINGS OF SMALL STS

### 6.1. Direct colourings.

- $v = 7$ ,  $\psi_{\max} = 3$   
System:  $\{0, 1, 2\}, \{0, 3, 4\}, \{0, 5, 6\}, \{1, 3, 5\}, \{1, 4, 6\}, \{2, 3, 6\}, \{2, 4, 5\}$   
Colouring:  $0 \mapsto 1, 1 \mapsto 1, 2 \mapsto 2, 3 \mapsto 3, 4 \mapsto 3, 5 \mapsto 1, 6 \mapsto 2$
- $v = 9$ ,  $\psi_{\max} = 4$   
System:  $\{0, 1, 2\}, \{0, 3, 6\}, \{0, 4, 8\}, \{0, 5, 7\}, \{1, 3, 8\}, \{1, 4, 7\}, \{1, 5, 6\}, \{2, 3, 7\}, \{2, 4, 6\}, \{2, 5, 8\}, \{3, 4, 5\}, \{6, 7, 8\}$   
Colouring:  $0 \mapsto 1, 1 \mapsto 2, 2 \mapsto 2, 3 \mapsto 1, 4 \mapsto 3, 5 \mapsto 4, 6 \mapsto 3, 7 \mapsto 4, 8 \mapsto 4$
- $v = 13$ ,  $\psi_{\max} = 5$   
System:  $\{0, 1, 4\}, \{1, 2, 5\}, \{2, 3, 6\}, \{3, 4, 7\}, \{4, 5, 8\}, \{5, 6, 9\}, \{6, 7, 10\}, \{7, 8, 11\}, \{8, 9, 12\}, \{9, 10, 0\}, \{10, 11, 1\}, \{11, 12, 2\}, \{12, 0, 3\}, \{0, 2, 7\}, \{1, 3, 8\}, \{2, 4, 9\}, \{3, 5, 10\}, \{4, 6, 11\}, \{5, 7, 12\}, \{6, 8, 0\}, \{7, 9, 1\}, \{8, 10, 2\}, \{9, 11, 3\}, \{10, 12, 4\}, \{11, 0, 5\}, \{12, 1, 6\}$   
Colouring:  $0 \mapsto 2, 1 \mapsto 4, 2 \mapsto 4, 3 \mapsto 4, 4 \mapsto 5, 5 \mapsto 3, 6 \mapsto 3, 7 \mapsto 3, 8 \mapsto 5, 9 \mapsto 2, 10 \mapsto 1, 11 \mapsto 4, 12 \mapsto 5$

- $v = 15$ ,  $\psi_{\max} = 5$   
 System:  $\{0, 1, 2\}$ ,  $\{0, 3, 7\}$ ,  $\{0, 4, 12\}$ ,  $\{0, 5, 8\}$ ,  $\{0, 6, 11\}$ ,  $\{0, 9, 13\}$ ,  $\{0, 10, 14\}$ ,  
 $\{3, 4, 5\}$ ,  $\{1, 4, 10\}$ ,  $\{1, 7, 13\}$ ,  $\{1, 9, 14\}$ ,  $\{1, 3, 9\}$ ,  $\{1, 8, 12\}$ ,  $\{1, 5, 6\}$ ,  $\{6, 7, 8\}$ ,  
 $\{2, 8, 14\}$ ,  $\{2, 6, 10\}$ ,  $\{2, 7, 9\}$ ,  $\{2, 5, 12\}$ ,  $\{2, 4, 11\}$ ,  $\{2, 3, 13\}$ ,  $\{9, 10, 11\}$ ,  
 $\{5, 9, 13\}$ ,  $\{3, 8, 11\}$ ,  $\{3, 10, 12\}$ ,  $\{4, 7, 14\}$ ,  $\{3, 14, 6\}$ ,  $\{4, 8, 9\}$ ,  $\{12, 13, 14\}$ ,  
 $\{6, 9, 12\}$ ,  $\{5, 9, 14\}$ ,  $\{4, 6, 13\}$ ,  $\{8, 10, 13\}$ ,  $\{5, 7, 10\}$ ,  $\{7, 11, 12\}$   
 Colouring:  $0 \mapsto 1$ ,  $1 \mapsto 2$ ,  $2 \mapsto 1$ ,  $3 \mapsto 3$ ,  $4 \mapsto 2$ ,  $5 \mapsto 4$ ,  $6 \mapsto 4$ ,  $7 \mapsto 3$ ,  $8 \mapsto 3$ ,  
 $9 \mapsto 4$ ,  $10 \mapsto 5$ ,  $11 \mapsto 4$ ,  $12 \mapsto 3$ ,  $13 \mapsto 1$ ,  $14 \mapsto 1$

**6.2. Computer constructions.** In each example below, a PTS is given from the proof of Theorem 4.2. A simple hill-climbing algorithm embeds each PTS in an STS( $v$ ). Each computation took a few seconds on a personal computer.

- $v = 19$ ,  $\psi_{\max} = 6$   
 Forcing PTS:  
 $\{0, 1, 2\}$ ,  $\{0, 3, 4\}$ ,  $\{1, 3, 5\}$ ,  $\{4, 5, 6\}$ ,  $\{0, 6, 7\}$ ,  $\{1, 6, 8\}$ ,  $\{7, 8, 9\}$ ,  $\{0, 9, 10\}$ ,  
 $\{1, 9, 11\}$ ,  $\{3, 10, 11\}$ ,  $\{0, 12, 13\}$ ,  $\{1, 12, 14\}$ ,  $\{5, 13, 14\}$ ,  $\{7, 12, 15\}$ ,  $\{10, 13, 15\}$   
 Embedding:  
 $\{2, 3, 13\}$ ,  $\{4, 9, 15\}$ ,  $\{3, 7, 16\}$ ,  $\{5, 7, 18\}$ ,  $\{1, 4, 10\}$ ,  $\{2, 12, 18\}$ ,  $\{15, 17, 18\}$ ,  
 $\{5, 8, 16\}$ ,  $\{3, 8, 18\}$ ,  $\{9, 14, 18\}$ ,  $\{11, 13, 16\}$ ,  $\{8, 11, 15\}$ ,  $\{1, 13, 18\}$ ,  $\{4, 7, 13\}$ ,  
 $\{4, 11, 18\}$ ,  $\{2, 4, 8\}$ ,  $\{3, 9, 17\}$ ,  $\{0, 16, 18\}$ ,  $\{1, 15, 16\}$ ,  $\{6, 10, 18\}$ ,  $\{3, 6, 12\}$ ,  
 $\{8, 10, 12\}$ ,  $\{2, 6, 15\}$ ,  $\{6, 11, 14\}$ ,  $\{2, 5, 9\}$ ,  $\{2, 7, 11\}$ ,  $\{9, 12, 16\}$ ,  $\{5, 10, 17\}$ ,  
 $\{8, 13, 17\}$ ,  $\{1, 7, 17\}$ ,  $\{0, 5, 15\}$ ,  $\{2, 10, 16\}$ ,  $\{6, 16, 17\}$ ,  $\{6, 9, 13\}$ ,  $\{0, 11, 17\}$ ,  
 $\{5, 11, 12\}$ ,  $\{2, 14, 17\}$ ,  $\{7, 10, 14\}$ ,  $\{4, 12, 17\}$ ,  $\{4, 14, 16\}$ ,  $\{0, 8, 14\}$ ,  $\{3, 14, 15\}$
- $v = 21$ ,  $\psi_{\max} = 7$   
 Forcing PTS:  
 $\{0, 1, 2\}$ ,  $\{0, 3, 4\}$ ,  $\{1, 3, 5\}$ ,  $\{4, 5, 6\}$ ,  $\{0, 6, 7\}$ ,  $\{1, 6, 8\}$ ,  $\{7, 8, 9\}$ ,  $\{0, 9, 10\}$ ,  
 $\{1, 9, 11\}$ ,  $\{3, 10, 11\}$ ,  $\{0, 12, 13\}$ ,  $\{1, 12, 14\}$ ,  $\{5, 13, 14\}$ ,  $\{7, 12, 15\}$ ,  $\{10, 13, 15\}$ ,  
 $\{14, 15, 16\}$ ,  $\{0, 16, 17\}$ ,  $\{2, 16, 18\}$ ,  $\{4, 17, 18\}$ ,  $\{7, 16, 19\}$ ,  $\{11, 17, 19\}$   
 Embedding:  
 $\{5, 7, 20\}$ ,  $\{0, 8, 18\}$ ,  $\{9, 13, 18\}$ ,  $\{9, 12, 16\}$ ,  $\{5, 15, 18\}$ ,  $\{4, 19, 20\}$ ,  $\{2, 6, 15\}$ ,  
 $\{1, 10, 18\}$ ,  $\{3, 6, 16\}$ ,  $\{9, 14, 17\}$ ,  $\{14, 18, 19\}$ ,  $\{8, 15, 17\}$ ,  $\{2, 10, 19\}$ ,  $\{6, 10, 17\}$ ,  
 $\{5, 8, 11\}$ ,  $\{11, 13, 16\}$ ,  $\{3, 15, 20\}$ ,  $\{2, 3, 12\}$ ,  $\{0, 11, 15\}$ ,  $\{3, 7, 18\}$ ,  $\{2, 5, 9\}$ ,  
 $\{4, 7, 13\}$ ,  $\{0, 14, 20\}$ ,  $\{6, 13, 19\}$ ,  $\{5, 10, 16\}$ ,  $\{3, 8, 14\}$ ,  $\{4, 11, 12\}$ ,  $\{5, 12, 17\}$ ,  
 $\{6, 9, 20\}$ ,  $\{6, 12, 18\}$ ,  $\{1, 15, 19\}$ ,  $\{4, 9, 15\}$ ,  $\{1, 7, 17\}$ ,  $\{2, 17, 20\}$ ,  $\{8, 16, 20\}$ ,  
 $\{4, 8, 10\}$ ,  $\{11, 18, 20\}$ ,  $\{1, 4, 16\}$ ,  $\{1, 13, 20\}$ ,  $\{2, 7, 11\}$ ,  $\{10, 12, 20\}$ ,  $\{3, 13, 17\}$ ,  
 $\{7, 10, 14\}$ ,  $\{0, 5, 19\}$ ,  $\{8, 12, 19\}$
- $v = 25$ ,  $\psi_{\max} = 8$   
 Forcing PTS:  
 $\{0, 1, 2\}$ ,  $\{0, 3, 4\}$ ,  $\{1, 3, 5\}$ ,  $\{4, 5, 6\}$ ,  $\{0, 6, 7\}$ ,  $\{1, 6, 8\}$ ,  $\{7, 8, 9\}$ ,  $\{0, 9, 10\}$ ,  
 $\{1, 9, 11\}$ ,  $\{3, 10, 11\}$ ,  $\{0, 12, 13\}$ ,  $\{1, 12, 14\}$ ,  $\{5, 13, 14\}$ ,  $\{7, 12, 15\}$ ,  $\{10, 13, 15\}$ ,  
 $\{14, 15, 16\}$ ,  $\{0, 16, 17\}$ ,  $\{2, 16, 18\}$ ,  $\{4, 17, 18\}$ ,  $\{7, 16, 19\}$ ,  $\{11, 17, 19\}$ ,  
 $\{18, 19, 20\}$ ,  $\{0, 20, 21\}$ ,  $\{2, 20, 22\}$ ,  $\{4, 21, 22\}$ ,  $\{6, 20, 23\}$ ,  $\{10, 21, 23\}$ ,  
 $\{12, 22, 23\}$   
 Embedding:  
 $\{0, 14, 19\}$ ,  $\{2, 3, 9\}$ ,  $\{3, 8, 13\}$ ,  $\{2, 4, 23\}$ ,  $\{5, 11, 16\}$ ,  $\{7, 11, 21\}$ ,  $\{5, 7, 22\}$ ,  
 $\{8, 12, 17\}$ ,  $\{8, 11, 23\}$ ,  $\{0, 15, 23\}$ ,  $\{3, 17, 24\}$ ,  $\{8, 19, 24\}$ ,  $\{6, 13, 17\}$ ,  $\{3, 15, 20\}$ ,  
 $\{1, 13, 22\}$ ,  $\{10, 17, 22\}$ ,  $\{9, 13, 20\}$ ,  $\{1, 10, 16\}$ ,  $\{9, 14, 18\}$ ,  $\{2, 7, 14\}$ ,  $\{5, 8, 20\}$ ,  
 $\{2, 5, 19\}$ ,  $\{8, 15, 18\}$ ,  $\{2, 8, 10\}$ ,  $\{15, 22, 24\}$ ,  $\{3, 16, 22\}$ ,  $\{2, 15, 17\}$ ,  $\{1, 18, 23\}$ ,  
 $\{10, 18, 24\}$ ,  $\{8, 16, 21\}$ ,  $\{16, 20, 24\}$ ,  $\{6, 12, 16\}$ ,  $\{3, 14, 21\}$ ,  $\{11, 13, 24\}$ ,  
 $\{3, 7, 23\}$ ,  $\{2, 12, 24\}$ ,  $\{14, 17, 20\}$ ,  $\{1, 15, 19\}$ ,  $\{11, 12, 20\}$ ,  $\{0, 5, 24\}$ ,

{14, 23, 24}, {5, 17, 23}, {6, 18, 22}, {13, 19, 23}, {2, 6, 11}, {9, 19, 22},  
 {4, 10, 19}, {1, 4, 20}, {7, 10, 20}, {12, 19, 21}, {5, 9, 15}, {4, 9, 12}, {6, 15, 21},  
 {4, 7, 24}, {1, 7, 17}, {4, 8, 14}, {1, 21, 24}, {5, 10, 12}, {3, 12, 18}, {0, 8, 22},  
 {4, 11, 15}, {3, 6, 19}, {4, 13, 16}, {6, 9, 24}, {9, 16, 23}, {5, 18, 21}, {9, 17, 21},  
 {0, 11, 18}, {2, 13, 21}, {7, 13, 18}, {6, 10, 14}, {11, 14, 22}

**6.3. A recursive construction.** In some cases, recursive constructions of Steiner triple systems are amenable to complete colourings. We give one illustration of this.

**Lemma 6.1.** *Suppose there exists an STS( $u$ ),  $(U, \mathcal{A})$ , admitting a complete  $k$ -colouring with colour class sizes  $w_1, w_2, \dots, w_k$ . Let  $\sigma \in \mathcal{S}_k$  be some permutation. Suppose there exists a Latin square  $L$  of side  $n$  with row and column-disjoint  $w_i \times w_{\sigma(i)}$  sub-rectangles  $R_i$ , each of which contains the entries  $\{e_1, \dots, e_k\}$ . Then there exists an STS( $3u$ ) admitting a complete  $2k$ -colouring.*

*Proof.* We apply the standard tripling construction for Steiner triple systems. Let  $V = U \times \{1, 2, 3\}$ , and let  $U \times \{1\}$ ,  $U \times \{2\}$ ,  $U \times \{3\}$  index the rows, columns, and entries of  $L$ , respectively. Define a set of blocks  $\mathcal{B}$  on  $V$  by including  $A \times \{i\} \in \mathcal{B}$  for every  $A \in \mathcal{A}$ . In addition, put  $\{x, y, L(x, y)\} \in \mathcal{B}$  for every  $x \in U \times \{1\}$  and  $y \in U \times \{2\}$ . So  $(V, \mathcal{B})$  is an STS( $3u$ ), and we now describe a colouring of it. Colour the points in  $U \times \{1\}$  with a complete  $k$ -colouring of  $\{A \times \{1\} : A \in \mathcal{A}\}$ , and such that rows in subrectangle  $R_i$  get colour  $i$ . Likewise, colour  $U \times \{2\}$  so that columns of  $R_i$  receive colour  $i$ . Every pair of colours  $\{i, j\}$ ,  $1 \leq i < j \leq k$ , is now covered (at least twice). Colour the points in  $U \times \{3\}$  with a complete  $k$ -colouring of  $\{A \times \{3\} : A \in \mathcal{A}\}$ , using colours  $\{k+1, \dots, 2k\}$ . This covers pairs of colours  $\{k+i, k+j\}$ , where  $1 \leq i < j \leq k$ . We may arrange this latter colouring so that entry  $e_j$  receives colour  $k+j$  for  $j = 1, \dots, k$ . By hypothesis, every pair of colours  $\{i, k+j\}$ , with  $i, j \in \{1, \dots, k\}$ , is covered by  $\mathcal{B}$ .  $\square$

**Example 6.2.** *Using Lemma 6.1, we have  $\psi_{\max}(27) = 8$ . Use the complete 4-colouring of the STS(9) given earlier with the following Latin square. Observe that each of the four indicated rectangles contains entries 1,2,3,4.*

1	2	4	5	8	9	7	3	6
3	4	2	1	5	7	8	6	9
9	5	1	2	4	8	6	7	3
2	1	3	4	9	6	5	8	7
5	3	7	6	1	2	4	9	8
8	7	6	9	3	4	2	1	5
7	8	9	3	6	5	1	2	4
6	9	5	8	7	1	3	4	2
4	6	8	7	2	3	9	5	1

## ACKNOWLEDGEMENT

We would like to thank the anonymous referee, whose careful reading of this manuscript helped with the correctness and presentation of various results.

## REFERENCES

- [1] D. Bryant and D. Horsley, *A proof of Lindner's conjecture on embeddings of partial Steiner triple systems*, preprint.
- [2] N.-P. Chiang, *The achromatic numbers of some uniform hypergraphs*, Congr. Numer. **100** (1994), 245–250.
- [3] J. Nešetřil, K. T. Phelps, and V. Rödl, *On the achromatic number of simple hypergraphs*, Ars Combin. **16** (1983), 95–102.
- [4] L. Teirlinck, *On projective and affine hyperplanes*, J. Combin. Theory Ser. A **28** (1980), 290–306.

DEPARTMENT OF MATHEMATICS AND STATISTICS, UNIVERSITY OF VICTORIA,  
VICTORIA, BC, CANADA V8W 3P4  
*E-mail address:* dukes@math.uvic.ca

*E-mail address:* gmacgill@math.uvic.ca

*E-mail address:* partonk@math.uvic.ca